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Tracker Week

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2. Status 3. Plans Principle Barrel-endcap geometry Expected stub rates The procedure



\rightarrow Tracking trigger using associative memories (AM) in 1 slide:



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\rightarrow Some of the main difficulties:



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\rightarrow **Barrel+Endcap geometry:**



 \rightarrow For the moment we consider only the barrel part. Endcap is clearly another story (much higher rates).

 \rightarrow Barrel = 6 layers of stacked modules: 3PS (*pixel/strips*) and 3SS (*strips/strips*).

 \rightarrow In this presentation we looked at the **4 outermost layers** (*but we used the* **PS** *layer as an* **SS** *one*). It is planned to use the pixel info at some point, but we keep it simple for the moment....

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\rightarrow Data reduction via stubs:





 \rightarrow Using stubs instead of clusters reduces data rate by ~1 order of magnitude (*in 200PU* events)

 \rightarrow This rate reduction could be further enhanced by cuts on stub width (SW) and cluster width (CW). Values of these cuts depends on the Pt threshold you're looking for.



 \rightarrow CW and SW cuts can be applied online at the hardware level . Therefore a rough Pt selection can be applied before sending the info to L1 system.

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\rightarrow <u>Relation between stub width and Pt of the particle:</u>

 \rightarrow In our track trigger project, the Pt threshold is currently set to 2 GeV/c



 \rightarrow From these plots, one sees that an SW cut of 2/2.5 could significantly reduce the stub rate, with a relatively small effect on good data.

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\rightarrow Data loss estimation when cutting on stub width:



C C	Good data losses and raw rate with SW<3									
Layer	5	6	7	8	9	10				
Proportion of lost stubs (in %)	7,5	6,8	19,9	2,5	7,8	21,9				
Overall rate reducti	on 1,8	1,7	1,7	1,6	1,6	1,5				

- \rightarrow The gain on raw stub rate is significant.
- → The loss is layer dependent (SW increases with radius). Cut values can be optimized for each layer.
- \rightarrow We need to study how these cuts may affect the pattern matching process (for the moment we don't apply strict cuts during stub formation process).

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\rightarrow ROD stub rates using CW3 and SW2 cuts:



 \rightarrow For our study, we use stubs made with CW<7 and SW<5 (conservative approach).

 \rightarrow The cuts are applied a posteriori for the moment.

 \rightarrow Suppose we need 20 bits to transmit one stub to the AM board. The rates per RODs are:

Layer	5	6	7	8	9	10
Max rate (in GHz/ROD)	32,3	17,7	9,4	6,0	3,3	2,0
Rate with stub cuts (in GHz/ROD)	17,6	10,4	5,6	3,7	2,1	1,3

 \rightarrow Using only the 4 last layers, 10GHz/ROD seems sufficient in all cases (*ie w or w/o stub cuts*).



\rightarrow Pattern matching: the procedure:

- 1. Choose a tracker geometry.
- 2. Depending on the constraints, try to find the best configuration using *clean* MC events.
 - \rightarrow The best configuration is the most efficient set of pattern/sectors.
 - \rightarrow The constraints are:
 - The minimum pT of the track you want to trigger (constrains the sector size)
 - The maximum number of words you could feed into the AM at L1 rate (constrains the sector size)
 - The maximum number of patterns you could store into one AM chip (constrains the pattern granularity)
 - The maximum number of words per pattern in order to ensure a correct fit (constrains the pattern granularity)
 - The maximum number of AM chips you can afford (constrains the sector size)
 - ...

3. The best configuration should be reasonably robust against pileup (efficiency should be robust, not fake rate), so step 2 has to be done also with heavy pileup MC samples.

4. If the results are not satisfying, one should restart at step 1

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Sector used and superstrip definition Bank generation Pattern matching on simple events Pattern matching on PU events



→ Sector definition:



- \rightarrow We define sectors using 'brute force' method:
- 1. We choose a ladder in the innermost layer
- 2. We look where are going the 2GeV/c $\mu^{\text{+/-}}$ crossing this ladder
- \rightarrow No sector overlap in the innermost layers.

 \rightarrow Of course, if necessary, these sectors can be divided into subsectors using independent pattern banks





 \rightarrow We tested 2 different sector sizes (3/4 layers)

 \rightarrow Using the two innermost layers seems challenging (*prohibitive rates*)

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\rightarrow Data rate in the AM with the defined sectors

 \rightarrow Using ROD rates presented previously (*the ones with stub cuts*), one could extrapolate the input rates in the AM board input buses, for the different configuations

Layer	5	6	7	8	9	10
Input rate for 3 layers (in GHz)				7,4	10,4	9,1
Input rate for 4 layers (in GHz)			5,6	7,4	10,4	10,4

 \rightarrow This can be reduced using smaller sectors, but it gives an idea of the numbers one should expect.

 \rightarrow Is this realistic or not???

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→ Superstrip definition:





 \rightarrow A superstrip is simply a bunch of strips in one module of the tracking detector.

 \rightarrow The superstrip address is the info sent to the AM board. Is is coded on a certain number of bits, depending on the superstrip resolution.

 \rightarrow The encoding is divided into 4 parts, giving module and intra-module SS position in Z and ϕ direction (*R* is not necessary)

 \rightarrow We are not using pixel info yet, so our Z intramodule encoding is very basic for the moment.

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\rightarrow The baseline:

- \rightarrow The bank generation software process iteratively using a set 20M Pgun μ^+/μ^- with p_T bet. 2 and 100 GeV/c
- \rightarrow The coverage is defined as the proportion of tracks in the sector matching a pattern in the bank.
- \rightarrow The coverage you require is one of the parameters driving the size of the bank.



- \rightarrow The superstrip size (see *plots*) is significantly affecting the bank size.
- \rightarrow With 32 strips, for an equivalent coverage, 3.5 times less patterns than with 16 strips are necessary
- → With 4 layers, one needs 700000 patterns to get a 90% coverage with 32 strips (not enough stat with 16 strips)

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\rightarrow <u>Reducing bank size using variable resolution patterns:</u>

 \rightarrow We adapted the technique described in <u>ATL-UPGRADE-PROC-2011-004</u>. During the bank generation, we associate the patterns to lower-res patterns:



 \rightarrow Using this method, you just need to construct one high resolution bank (*this assume that you have enough statistic to do it...*). The low-res patterns and linkings are done on-the-fly.

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\rightarrow <u>Reducing bank size using variable resolution patterns:</u>

 \rightarrow At the end, we loop over the low-res patterns found, and for each of them we retrieve all the corresponding high-res patterns. We **OR** them and compare the result to the parent lo-res pattern:





 \rightarrow If the two high-res SS are used, this means that we can't take advantage of the higher resolution for this pattern in this layer. Therefore we will set a don't care bit (X) in the low-res pattern. Don't care because there is nothing to gain here...

 \rightarrow On the other hand, if we see that only one of the high-res SS is fired in the merging result, then we can use the higher resolution. We will set a position bit (0 or 1) for this layer in the low-res pattern.

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\rightarrow <u>Reducing bank size using variable resolution patterns:</u>

 \rightarrow In the **pattern bank**, you keep the low-res pattern, and you just add one ternary bit per layer (*this is possible with ternary AM*). This bit has the three following states: X (don't care), 0, and 1.



 \rightarrow So basically you have the information of two hi-res patterns with one low-res pattern, you reduce the size of your bank but keep the sensitivity.

 \rightarrow You can add up to 3 DC bits in the current AM chip (*is it a fixed number*??), so we produced pattern banks with 1,2, or 3 DC bits, starting from the bank previously shown.

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\rightarrow Effect of DC bits on the pattern bank size:



 \rightarrow 90% coverage stands for the 0DC bank, it gets higher by construction when one adds DC bits.

 \rightarrow The reduction of the bank size is significant, but the low res patterns might be a bit large, in particular in the low res case. This could affect the fake rate.

 \rightarrow One need to start from higher resolution patterns (eg 4 layers with 16/8 strips). Much larger samples have to be produced.

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\rightarrow Software emulation principle:

 \rightarrow In order to test the efficiency of the generated bank, a software emulation of the AM was written.

→ Works as the bank generator (*ie fully flexible*). Takes in input the **pattern bank** and the **stubs of a given event**.

 \rightarrow As in the AM, the pattern ID is done in one single pass. Output for each event is the list of patterns activated.



 \rightarrow CPU consumption of the pattern matching w.r.t. the bank size

 \rightarrow The algorithm developed is linear w.r.t. the bank size.

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→ Few definitions:



Fake rate = -	# of activated patterns containing NO primary track					
	# of activated patterns					



 \rightarrow We want a good efficiency for particle with Pt>2GeV/c, with the lowest possible fake rate, and keep the redundancy a low as possible.

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→ <u>Results on mu+/mu-:</u>

ightarrow Baseline result for 3 layer sector, using the classic bank, requiring exactly 3 stubs on the pattern



 \rightarrow Slow turn-on curve (we're aiming for 2GeV/c), but low fake rate and redundancy

 \rightarrow The slow turn on comes from tracks with missing stubs, and from the fact that it's difficult to populate the bank at low Pt.

 \rightarrow First point can be solved by adding missing stubs possibility, second point by adding DC bits.

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\rightarrow <u>Result for electrons (worst case for single particle)</u>:



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\rightarrow Conclusion on simple events:

- \rightarrow The addition of DC bits and missing stubs is significantly improving the turn on of efficiency curve.
- \rightarrow For muons (and also for pions), we can be within the requirements (*Pt>2GeV/c with >90% efficiency*).
- \rightarrow The convergence for electrons is also pretty good, but as expected a bit worse.

 \rightarrow The price to pay is a larger redundancy (*when adding DC bits*), and a larger fake rate (*when adding missing stubs*). Depending on the difficulty to sort this out at the trackfit stage, this could be or not a problem.

 \rightarrow Next step is to look at heavy pile up events (200PU).

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\rightarrow Rates of matched patterns with 200 PU events at 20MHz:

	N DC bits	0	1	2	3
Sector with 3 layers	Total pattern rate (in MHz)	11,2	12,4	22,2	52,4
	Good pattern rate (in MHz)	7,9	8,1	10,5	17,1
Sector with 4 layers	Total pattern rate (in MHz)	219,8	214,2	428,6	1073,2
	Good pattern rate (in MHz)	69,7	57,3	70,0	105,2

 \rightarrow The fake rate is clearly larger in PU events (*not really a surprise*).

 \rightarrow The efficiency looks small, but is mainly driven by poor efficiency on the huge amount of low Pt particles. At higher Pt the efficiency is close to 1.

 \rightarrow We reach 90% at ~15GeV/c. This is far from 2GeV/c, but this is a very first look, there is plenty of room for improvement (*stub cuts, higher resolution pattern bank,...*)

 \rightarrow The 2GeV/c is not out of reach at all...



 \rightarrow We have set up a flexible software for studying AM-based tracking in CMS future tracker. This software is not depending on the geometry (*dealing with SS addresses*), and **provides bank** generation (*w or w/o DC bits*) and pattern matching tools.

- \rightarrow Different parameters have been tested, along with different event samples:
 - We start to get a better picture of the task, in particular we have a rough idea of the data rates that the system will have to sustain (using stub rates from PU events).
 - We see the effect of adding missing stubs and DC bits abilities. This leads to promising results for single particles, but fake rates becomes dominant for high PU events. But there is plenty of room for improvement here.
 - The superstrip granularity is strongly affecting the pattern bank size. Going from 16 to 32 strips doesn't affect the efficiency, but increase fake rate by 50% (*for PU*). One has to find a balance between precision and size of the system. DC bits can provide that, but we still have to fully understand there impact.
 - Considering the stub rates on the four outermost layers, **including a fourth layer seems a good option**, in particular in view of the track fit which comes after the AM matching. OK, but we definitely need more events to generate the pattern bank.
 - This fourth layer is a PS one. We should try to find a way to add the pixel info without affecting too much the bank size.