

# An attempt to innovate the standard model of control systems

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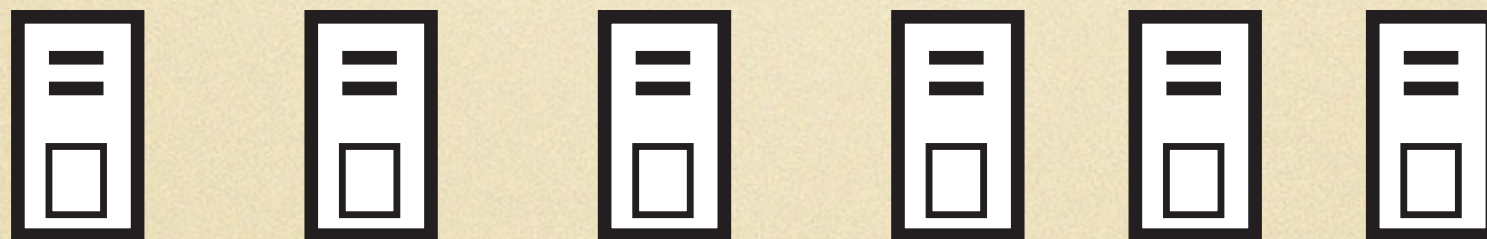
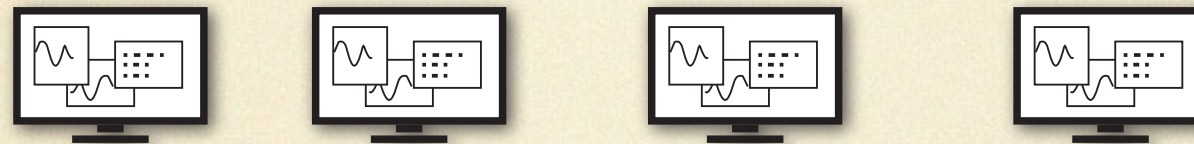
F. Zani (INFN Roma TV)



“The *standard model* consists of a local area network providing communication between front end microcomputers, connected to the accelerator, and workstations, providing the operator interface and computational support.”



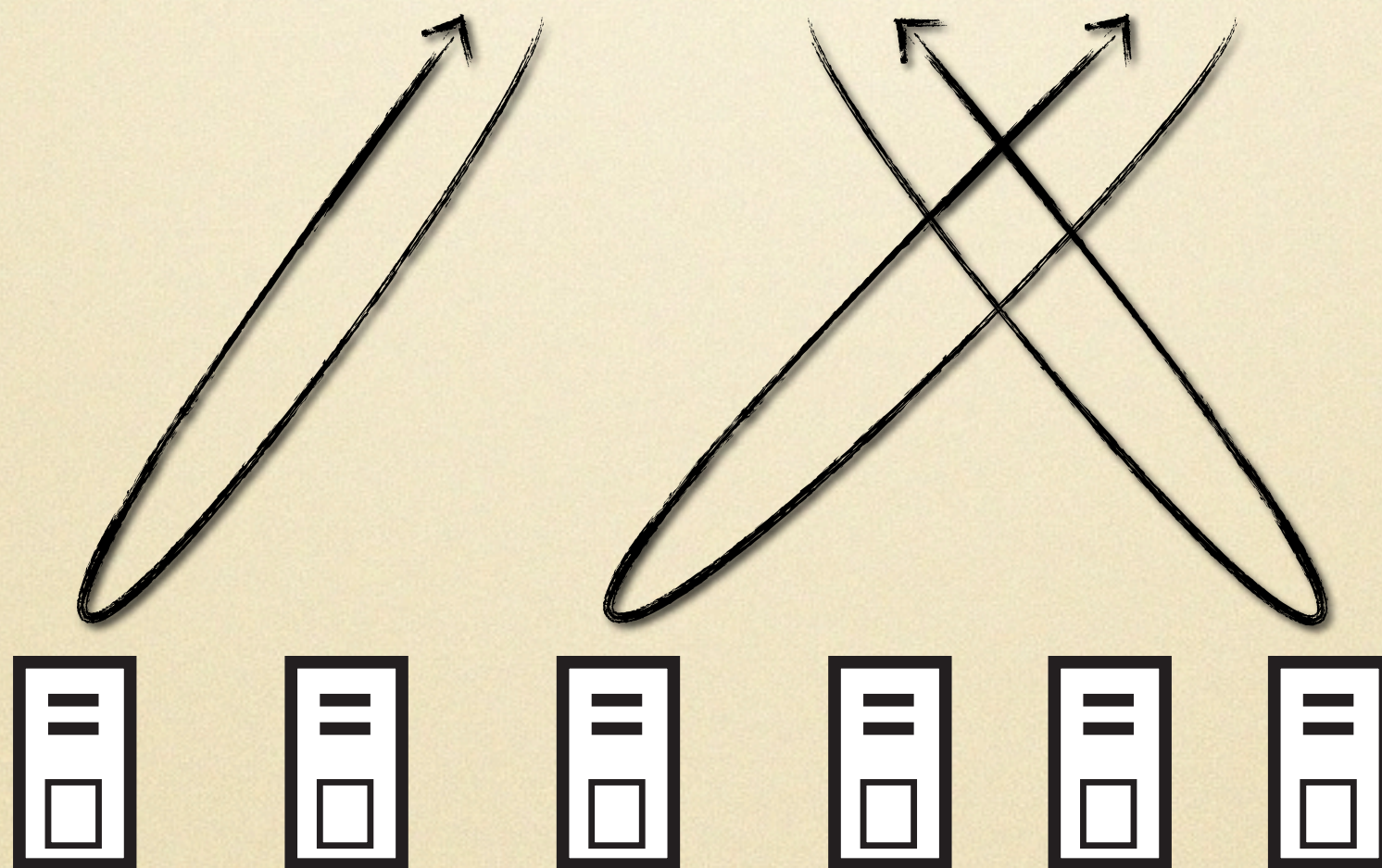
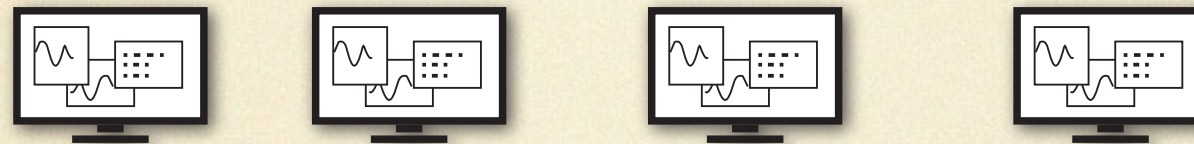
control room



accelerator



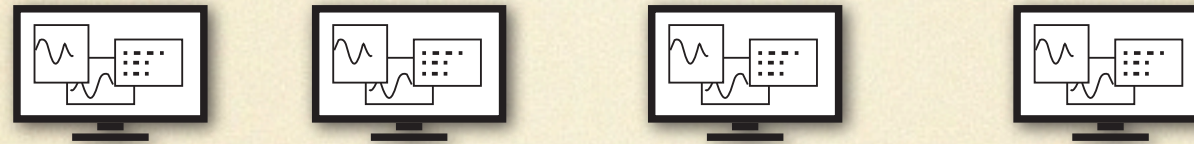
control room



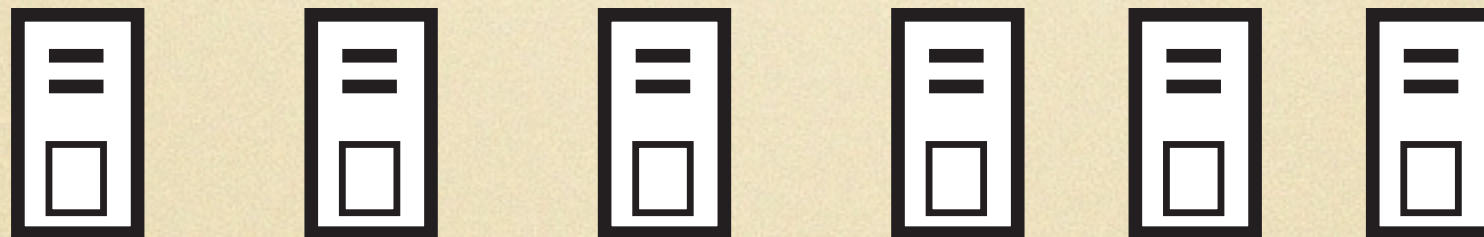
accelerator



control room



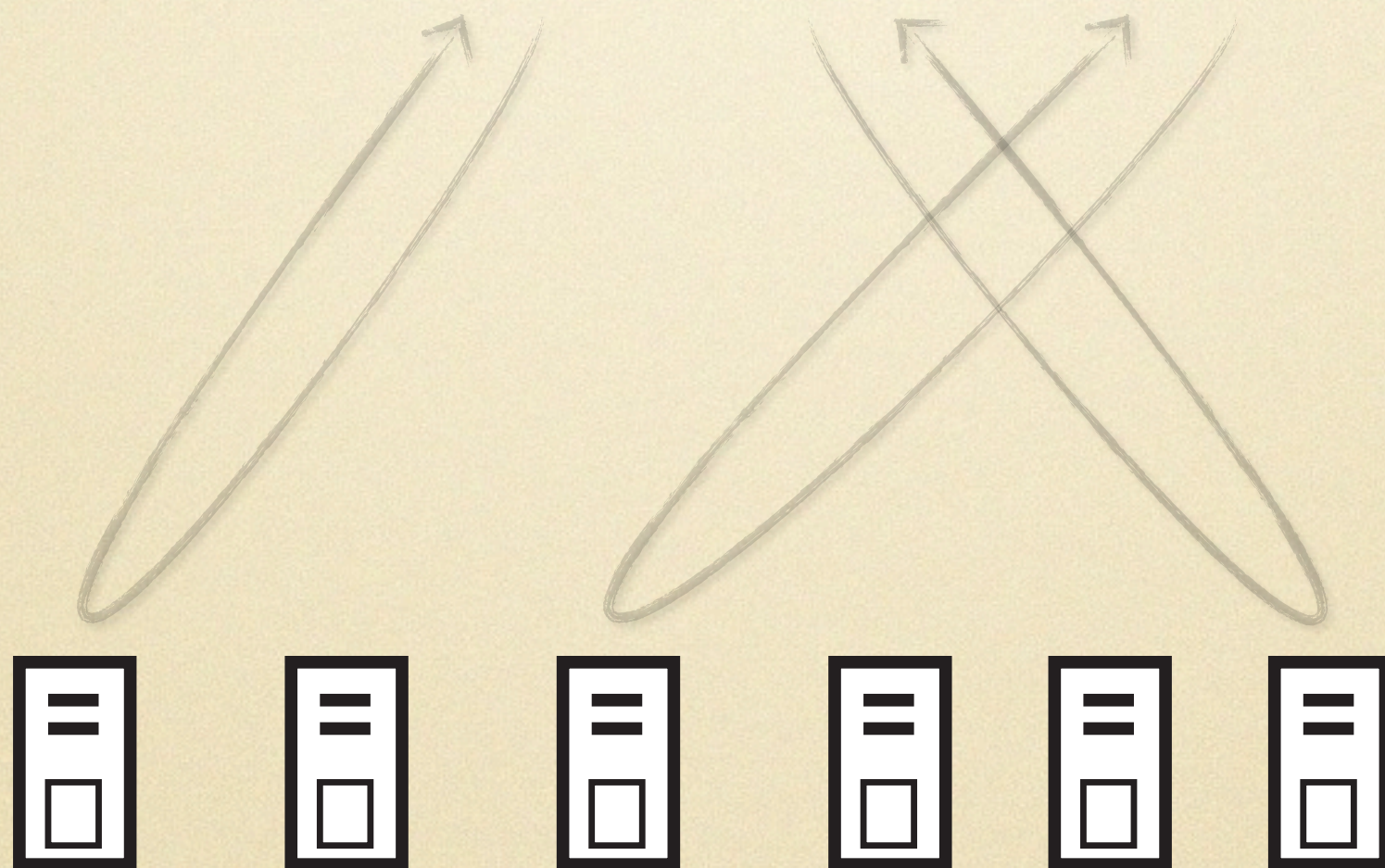
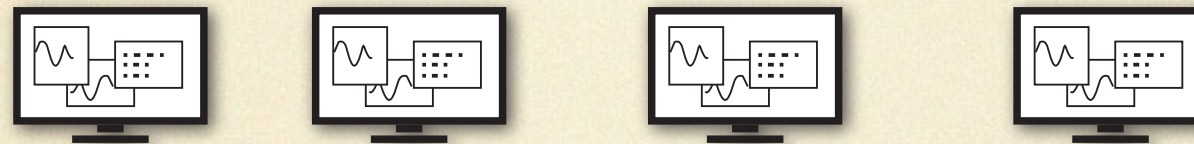
middle-layer



accelerator



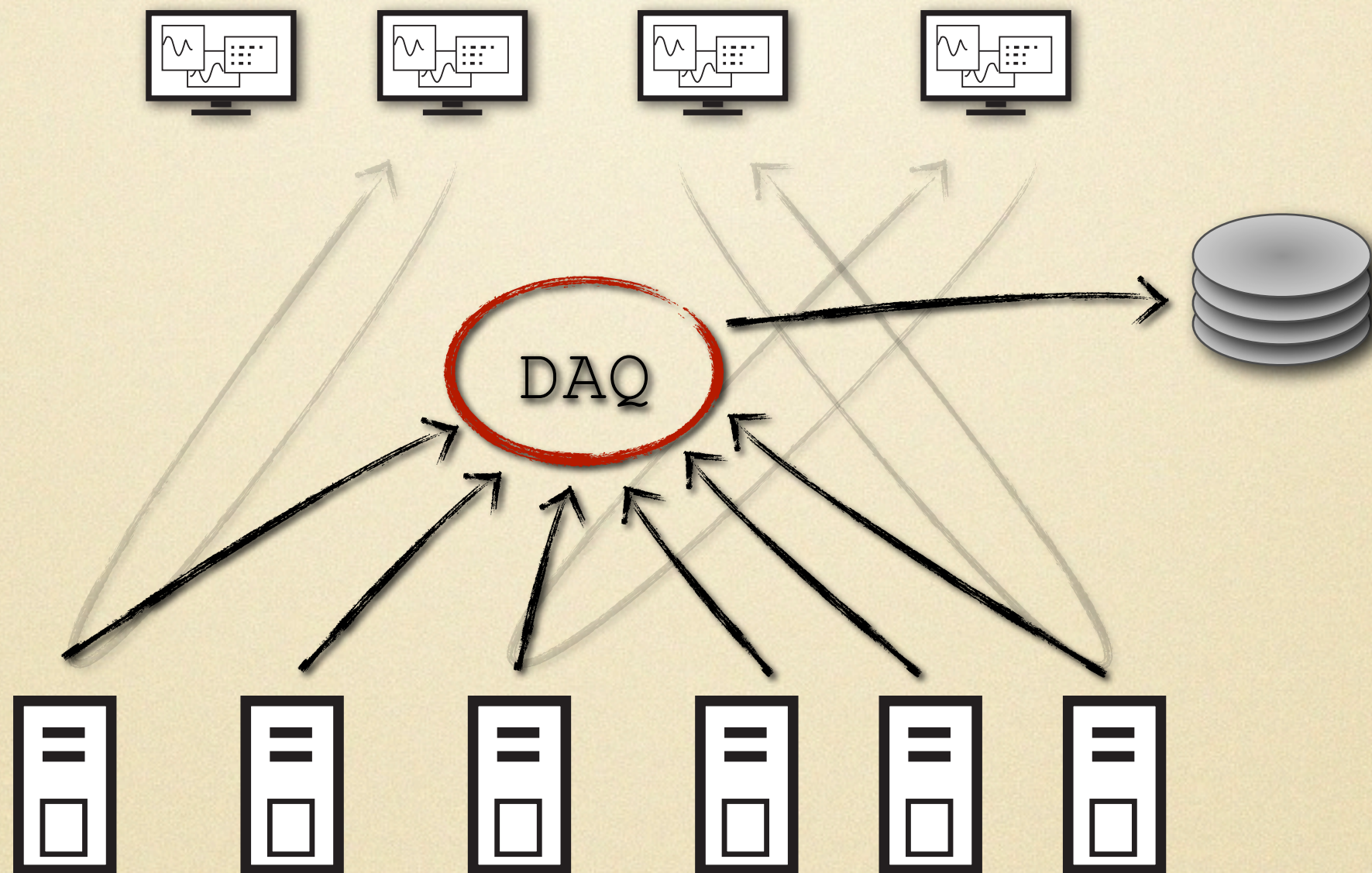
control room



accelerator



control room



accelerator



# the starting point

- goal: develop a new solution for a control system's DAQ
- use key / value db as alternative to RDBMS
  - fast, scalable, distributed storage, low-cost servers



the next step



# the next step

- extended goal: *key/value* db looks great, can we use it for live data ?
- no, data retrieving too slow

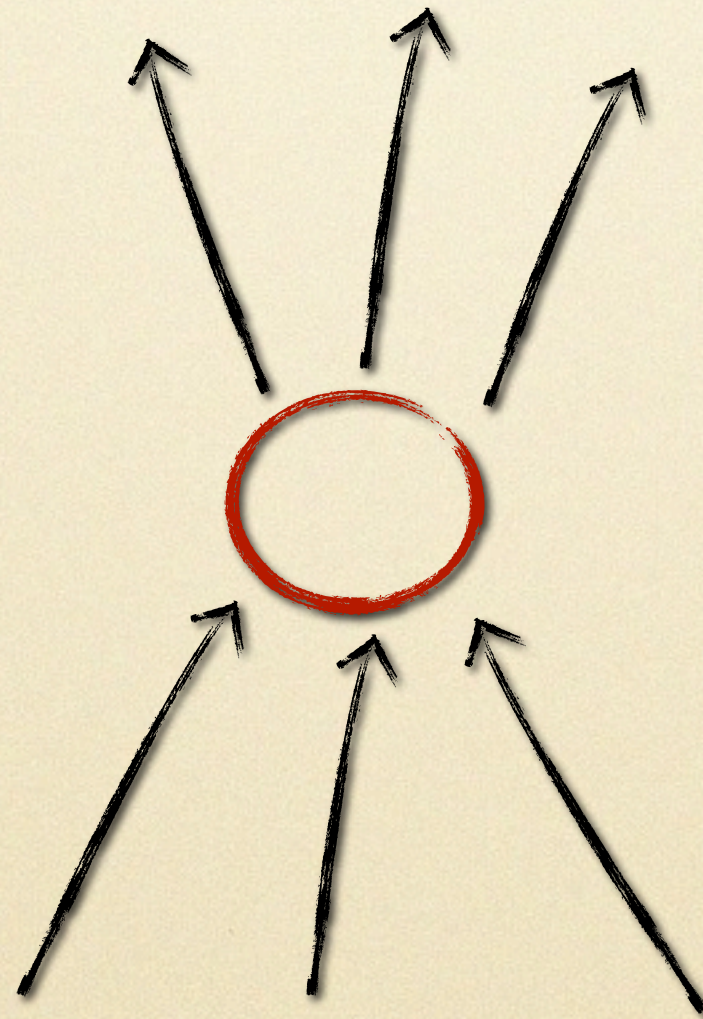


# the next step

- **extended goal:** *key/value* db looks great, can we use it for live data ?
  - no, data retrieving too slow
- use distributed caching instead
  - same topology, same data structure, similar scalability



control room



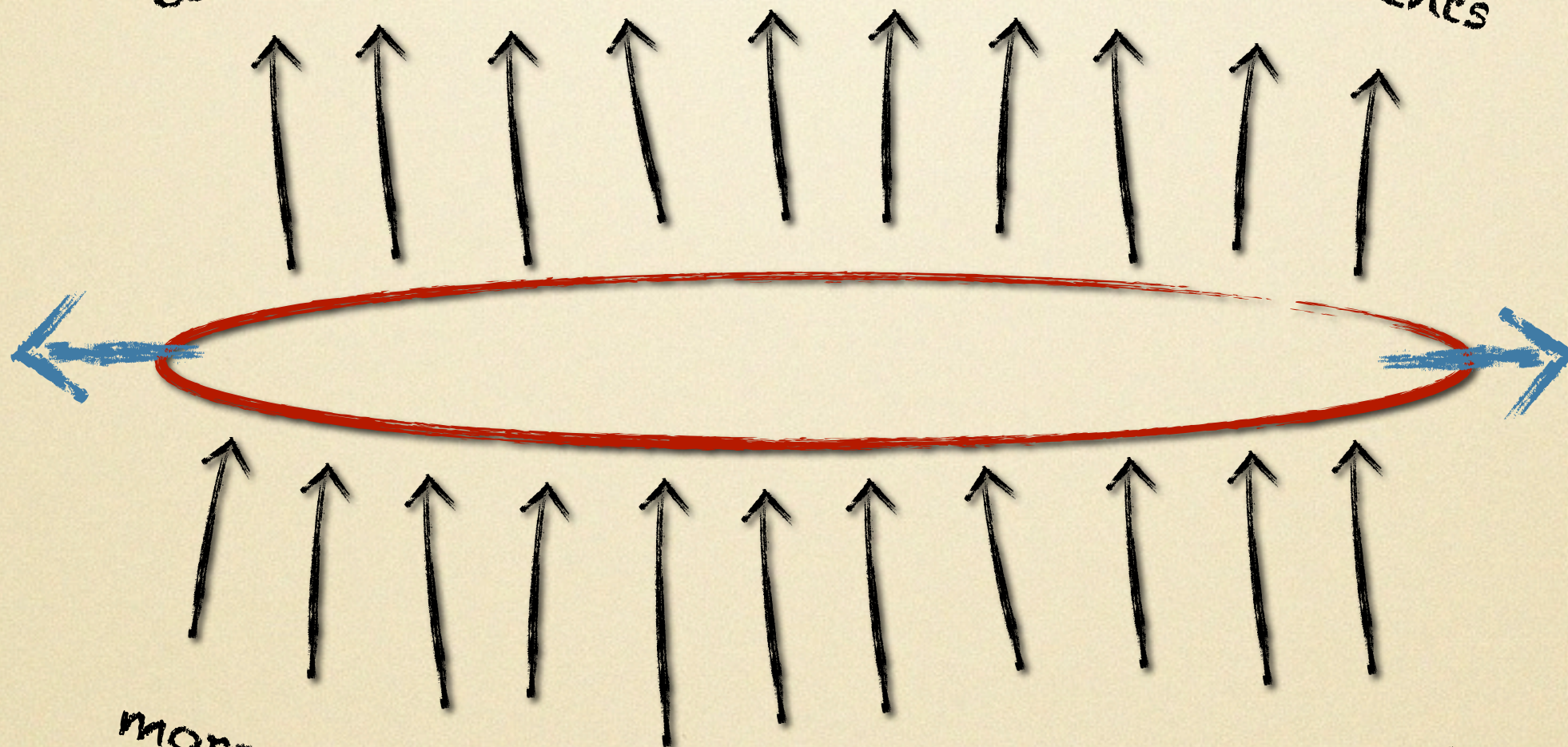
accelerator



control room

more  
clients

more  
clients

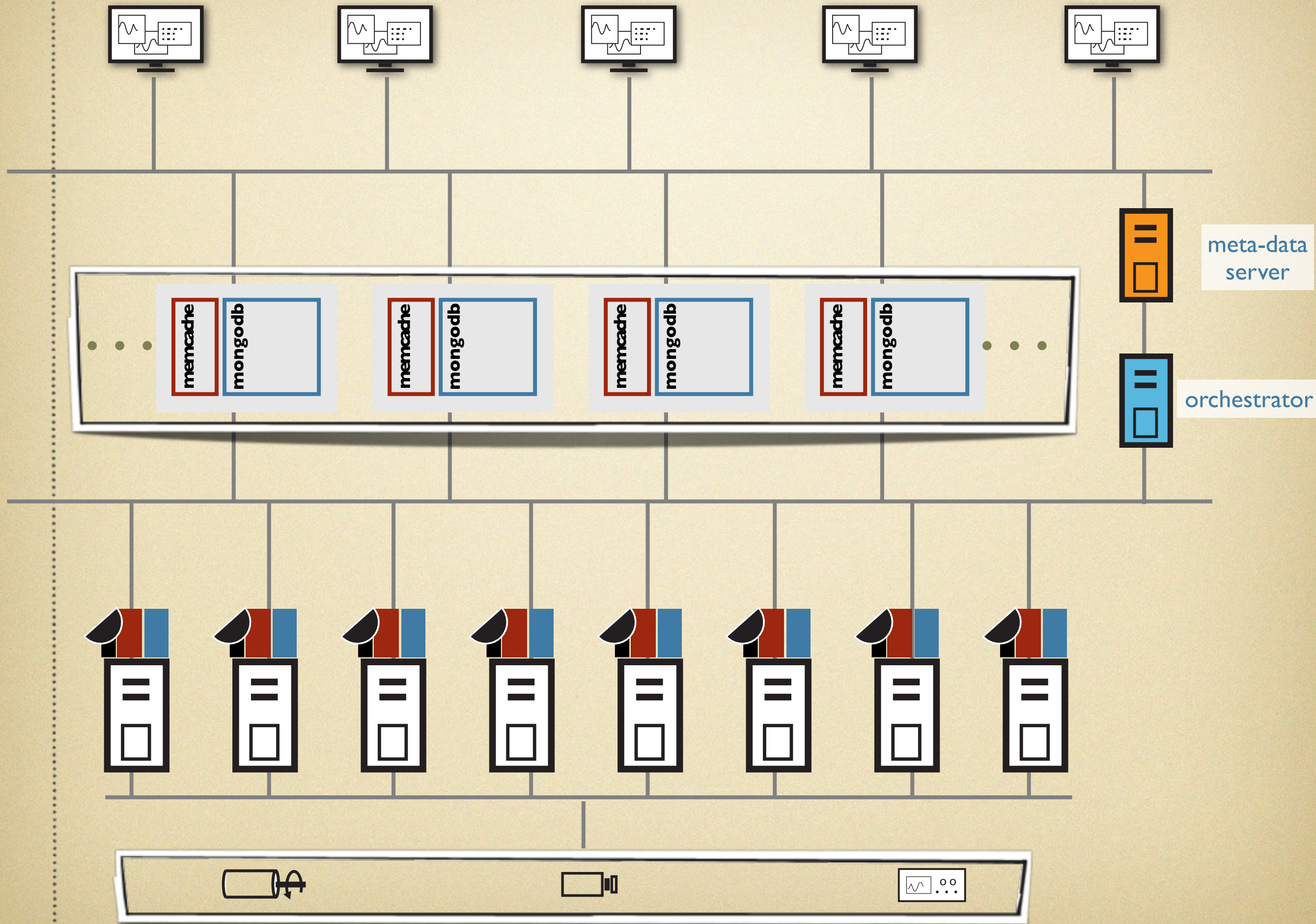


more  
devices

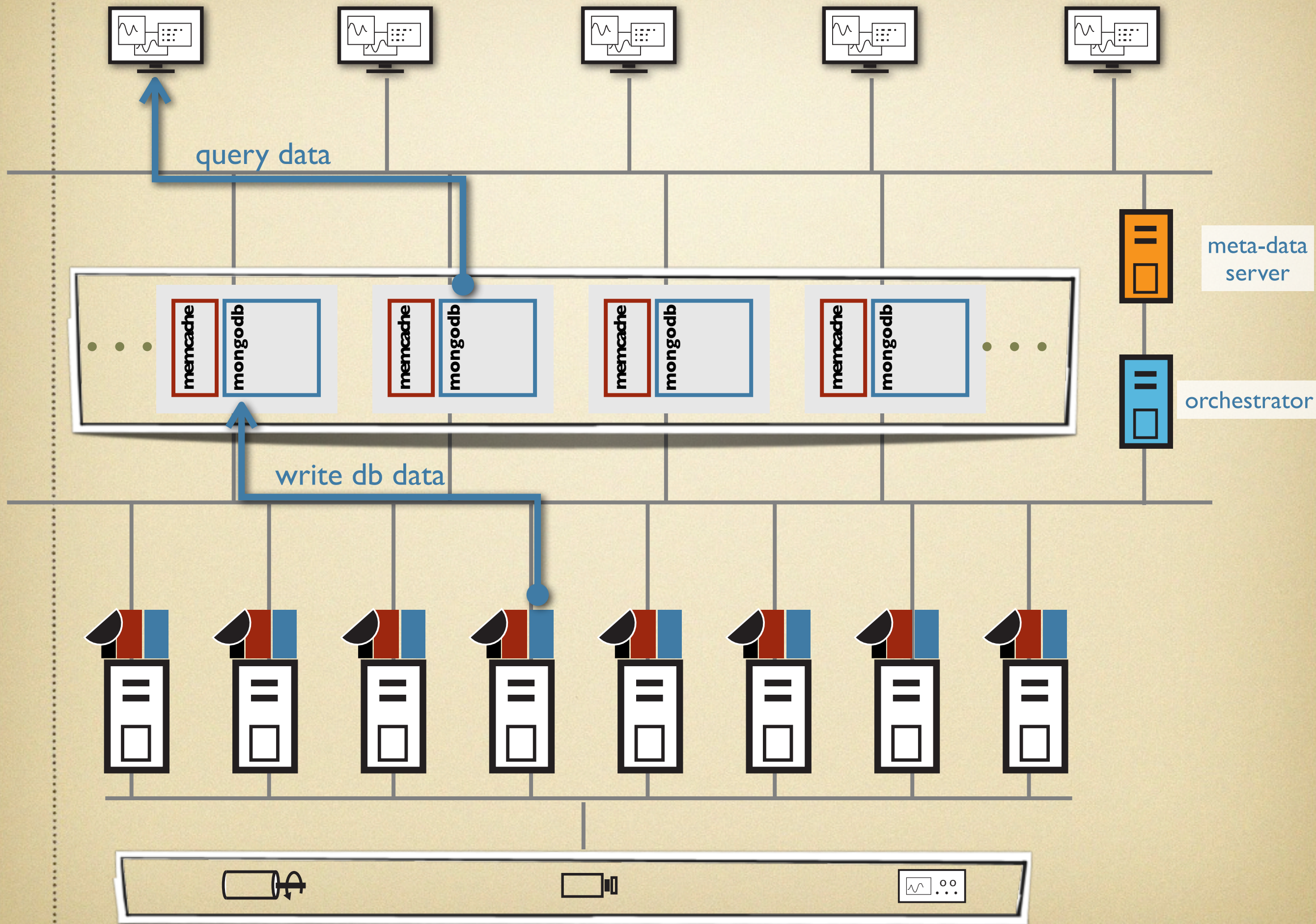
more  
devices

accelerator

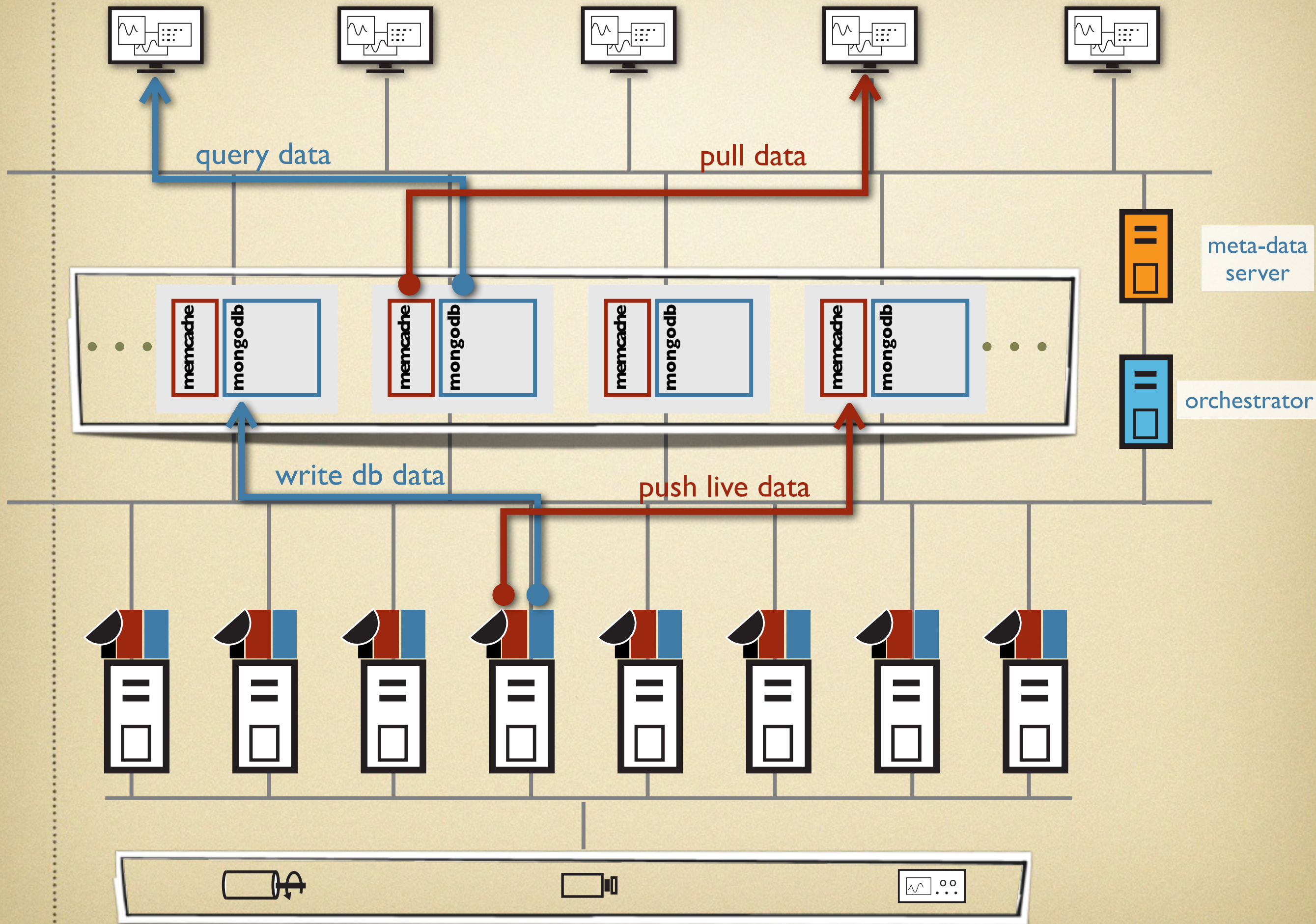




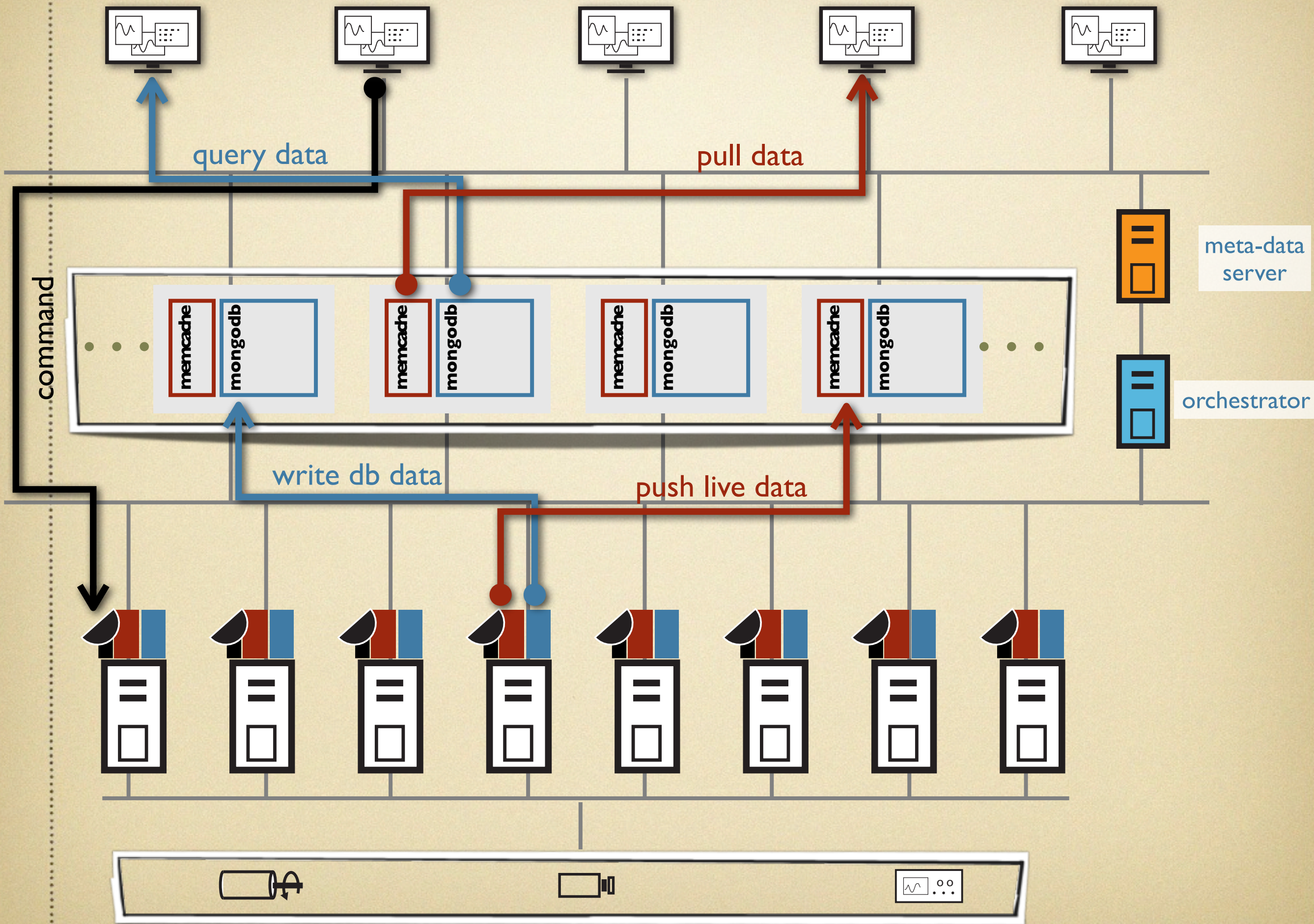






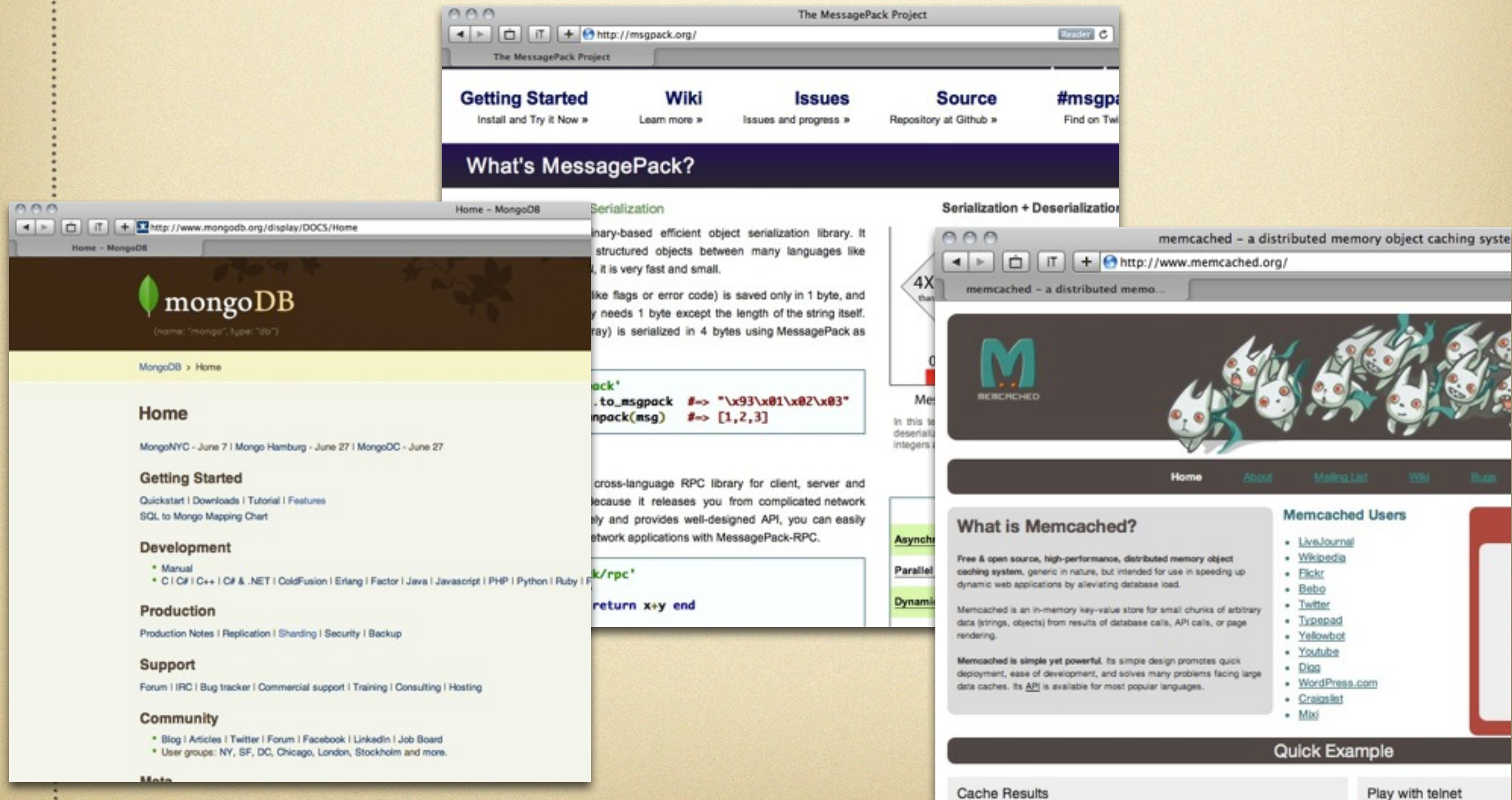






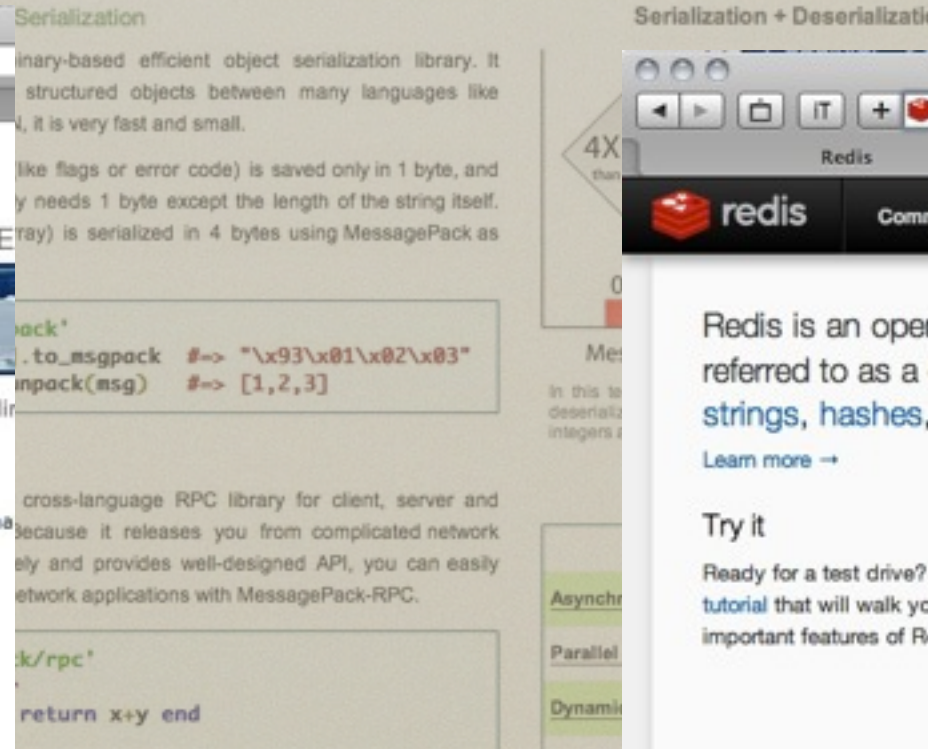
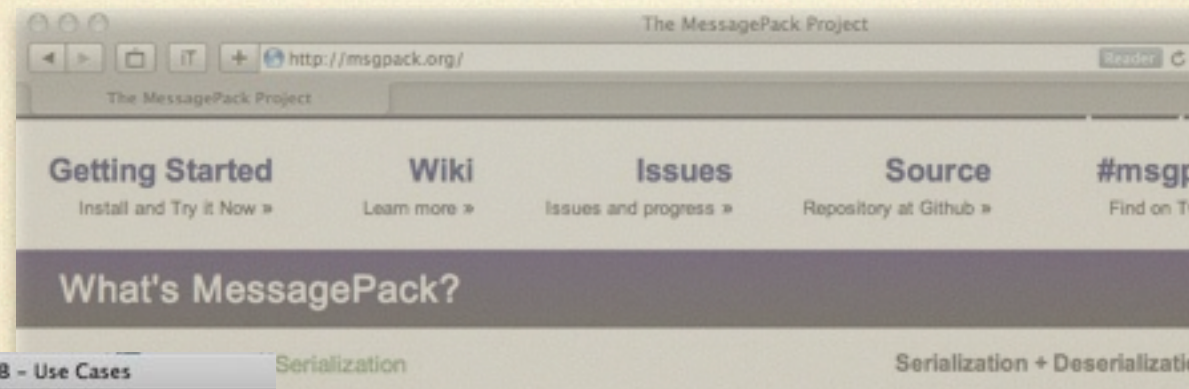


# core services candidates



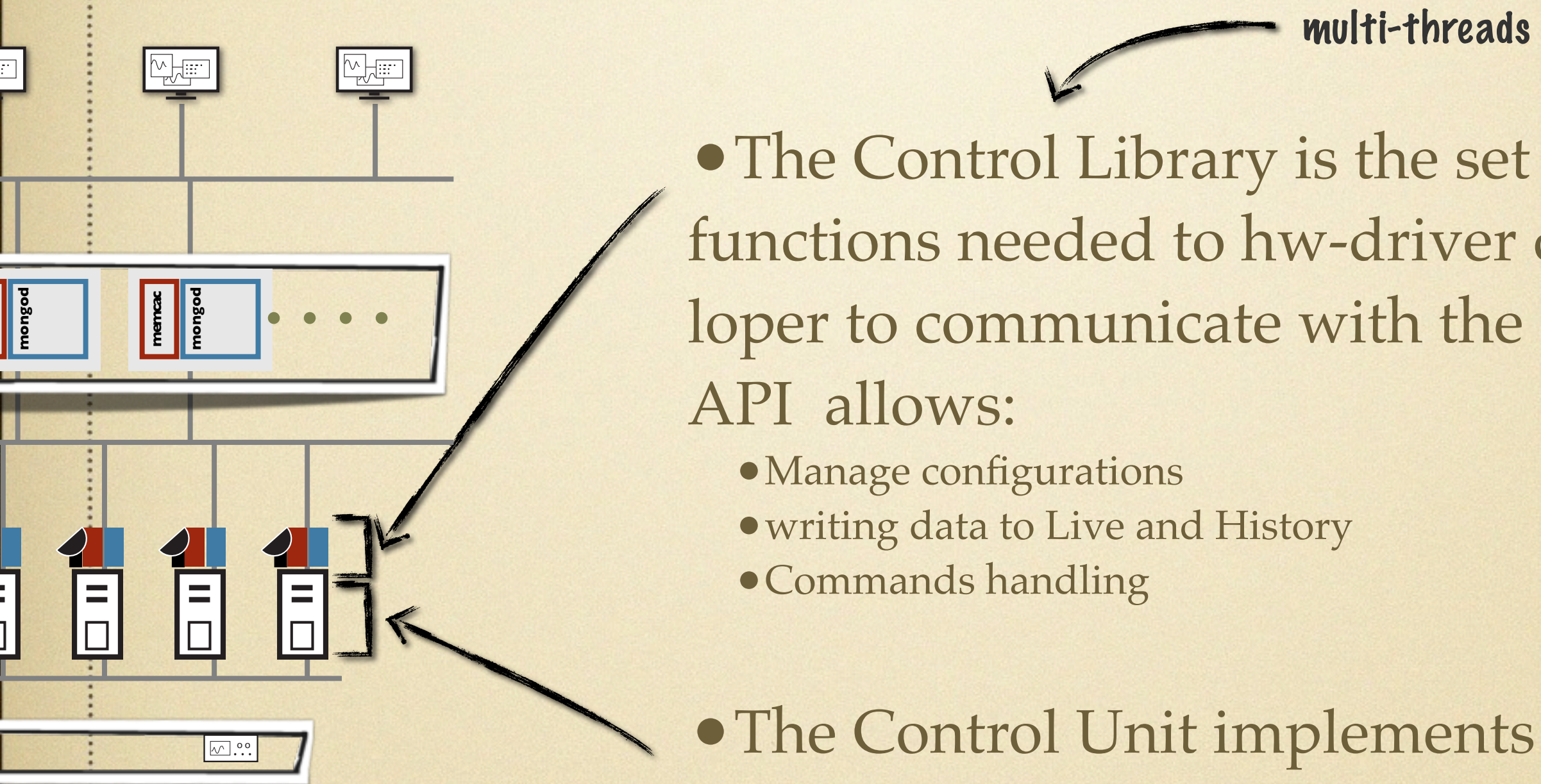


# core services alternatives





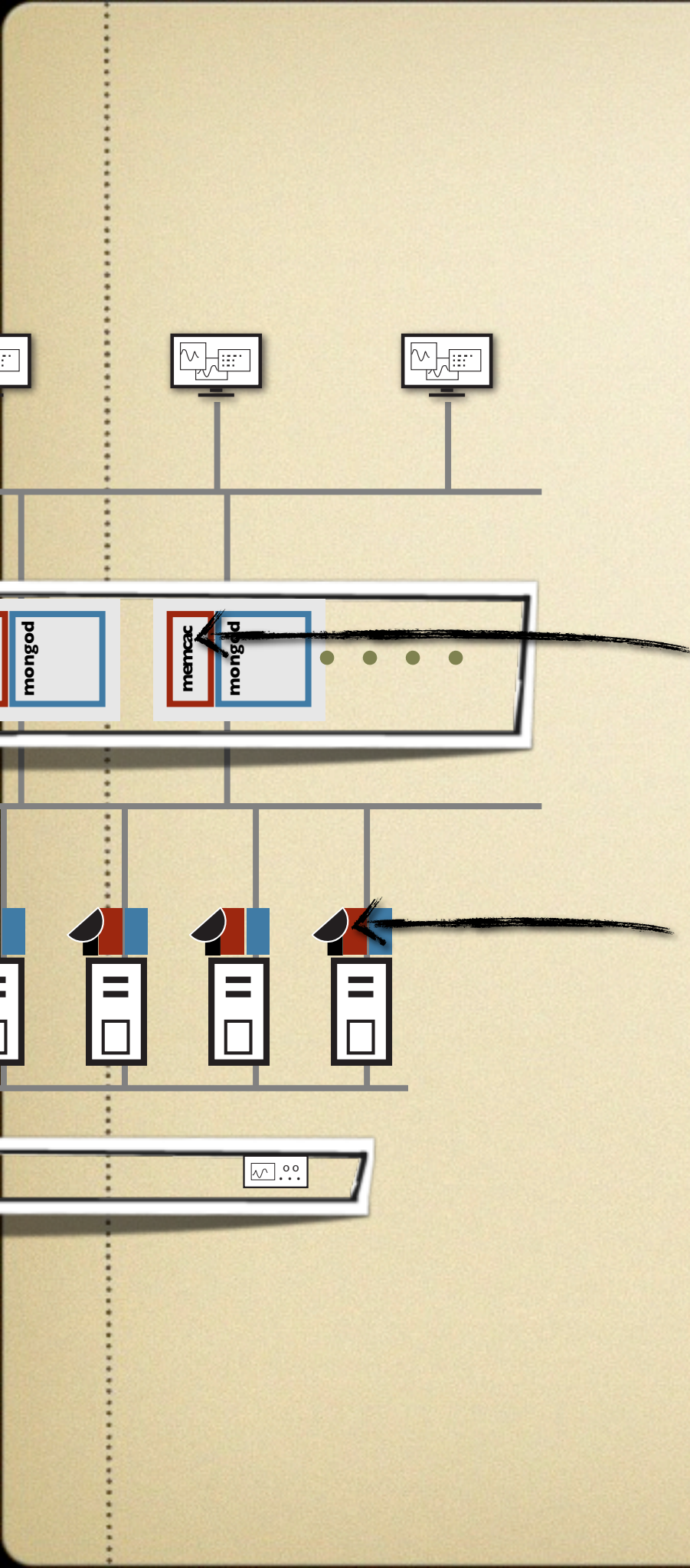
# Control Library e Control Unit





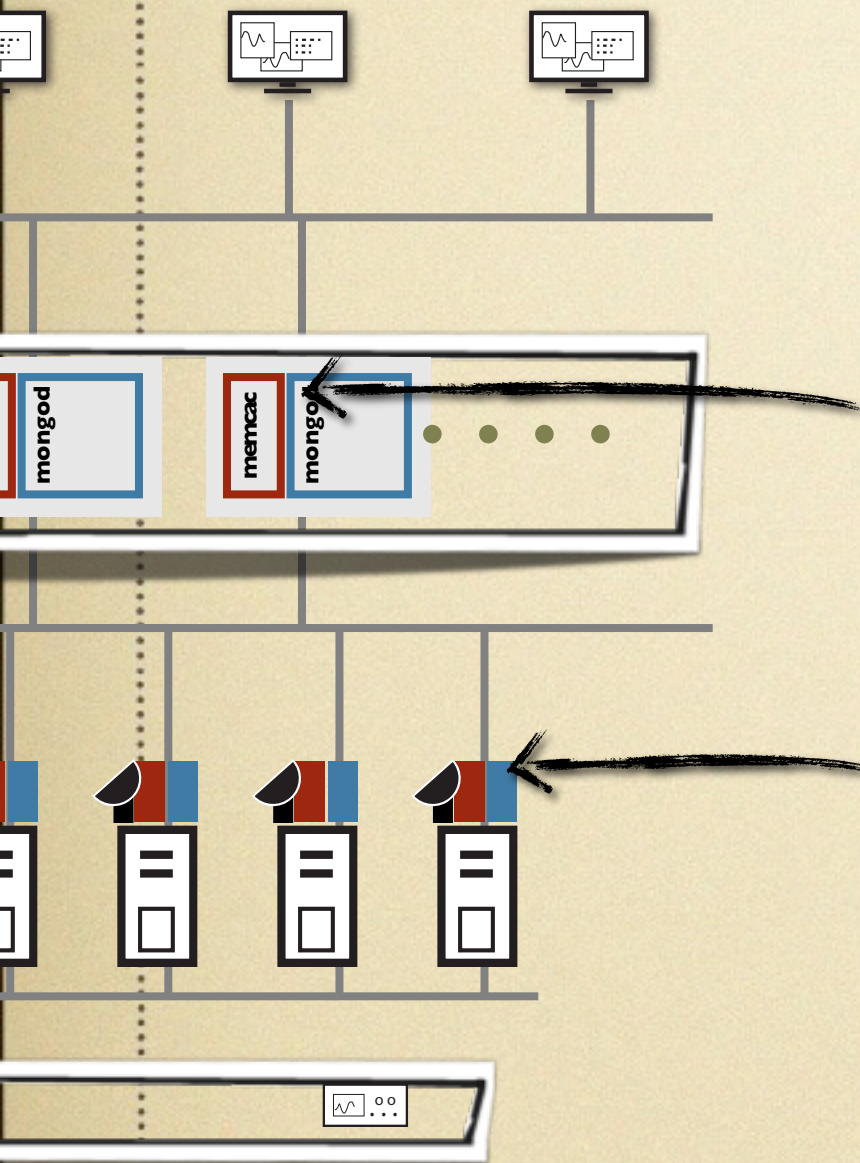
# Live data

- Allows high-performance caching of data produced by all components managed by CS.
- one *key* per data (a single “container” continuously updated)
- dynamical *keys* re-distribution allows automatic failover by distributing to other server the load of failed one.
- Scalability is also guaranteed by the same feature





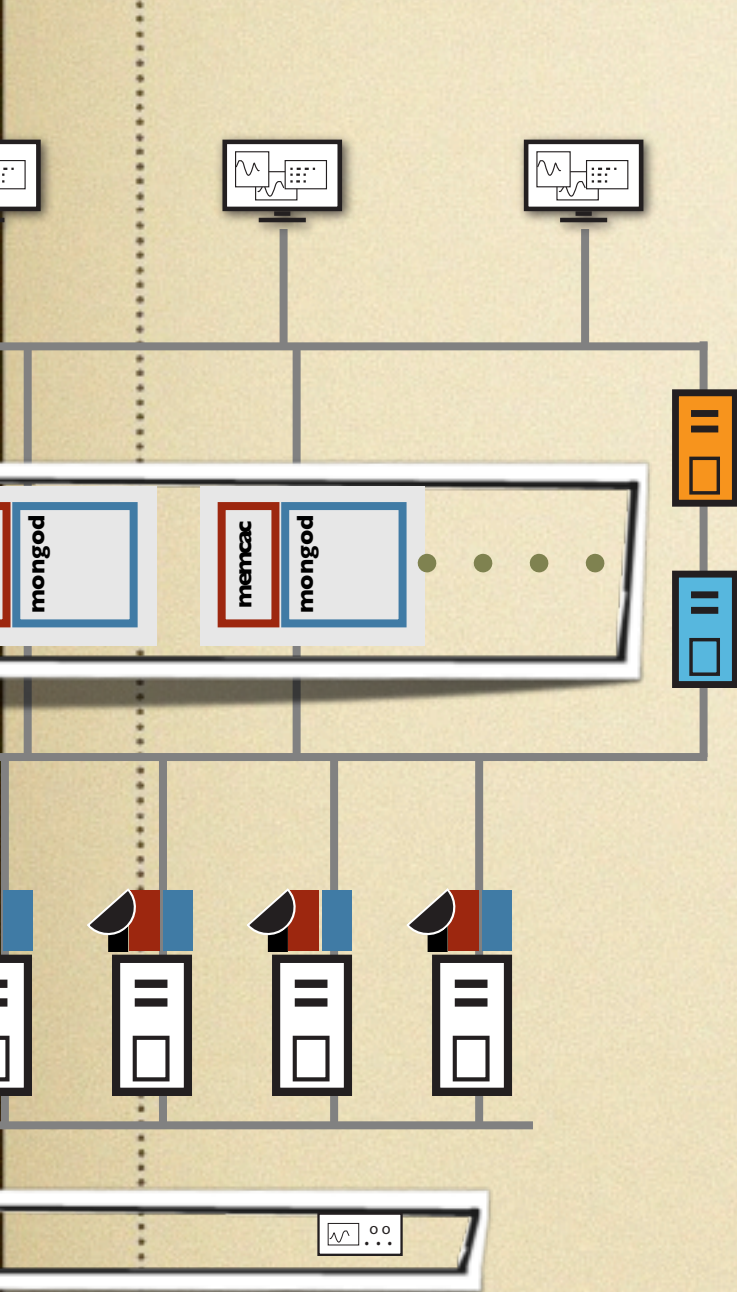
# History data



- *key/value* non-relational data-base
- scalability and load balancing by *sharding*
- fast record writing (simpler structure because it has no tables)
- fast queries on primary keys
- (fast) parallel search on cluster nodes



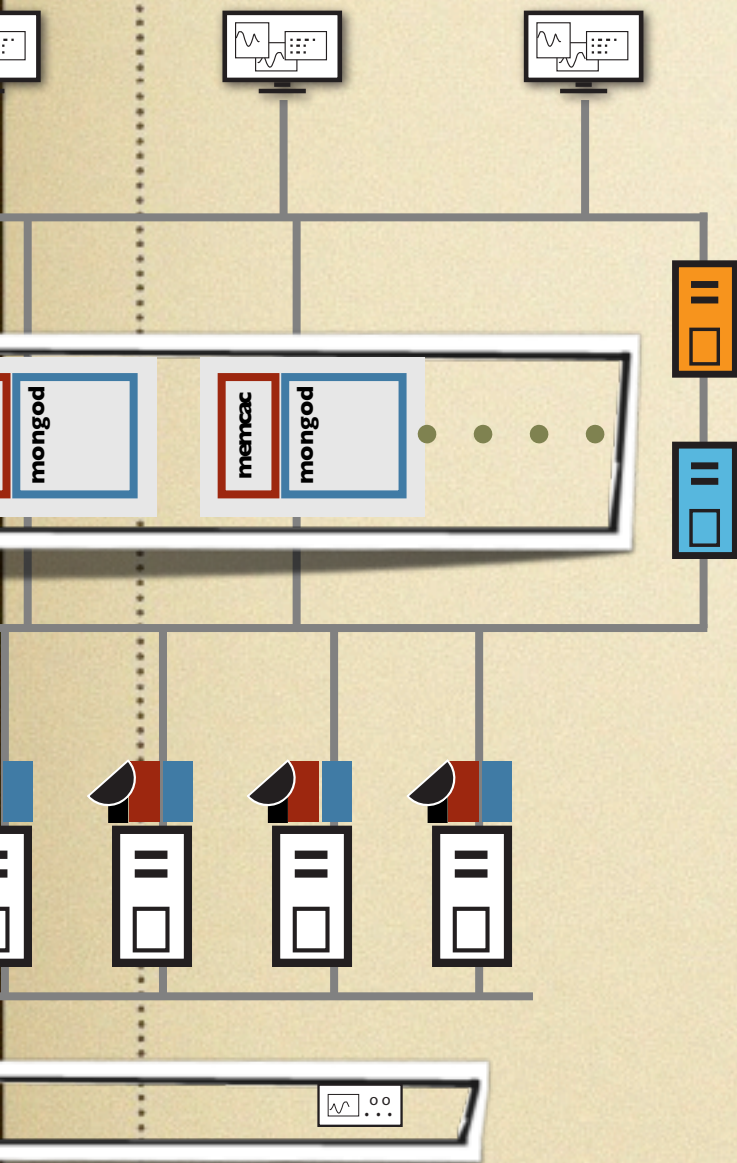
# Metadata Server



- CU configuration manager  
(e.g. managing of push data rate)
- Semantic of data (e.g. db records structure)
- Command's list and semantic
- Naming service



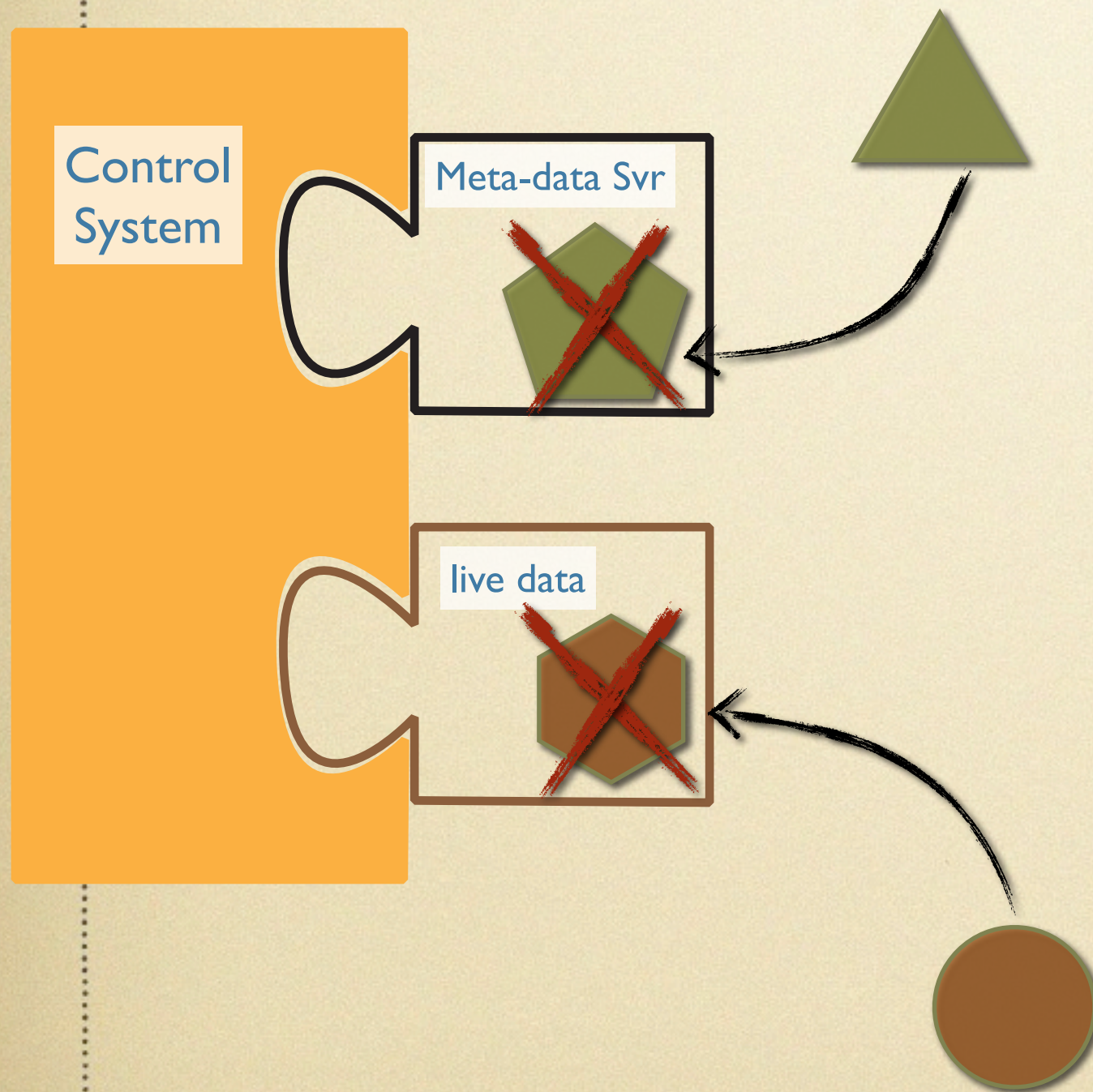
# Orchestrator



- Provides middle-layer services, e.g. locking of CUs to prevent command conflicts
- multi-CUs commands, e.g.
  - global set-points save/restore
  - software feedback
  - on-line measurements
  - ...



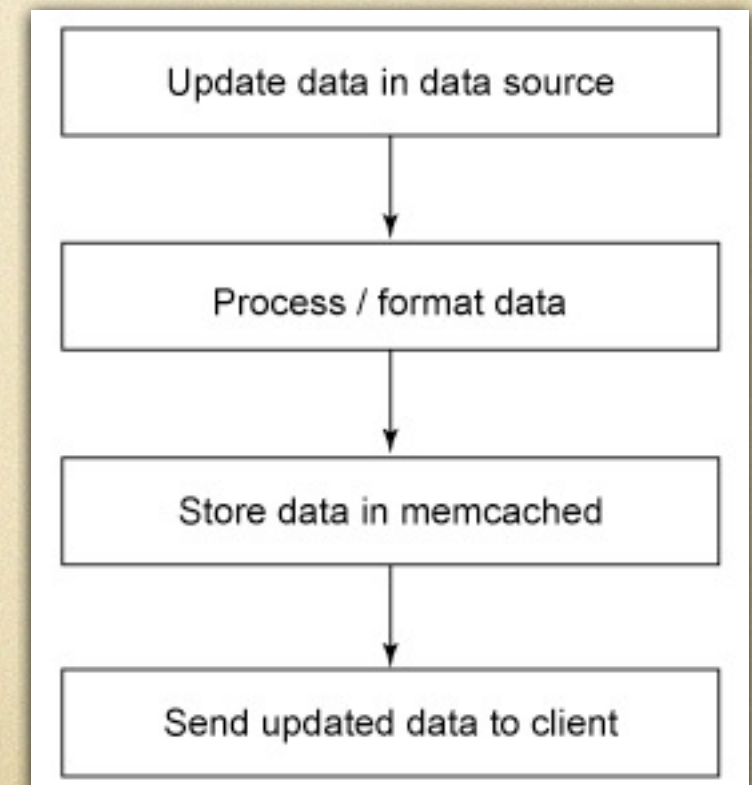
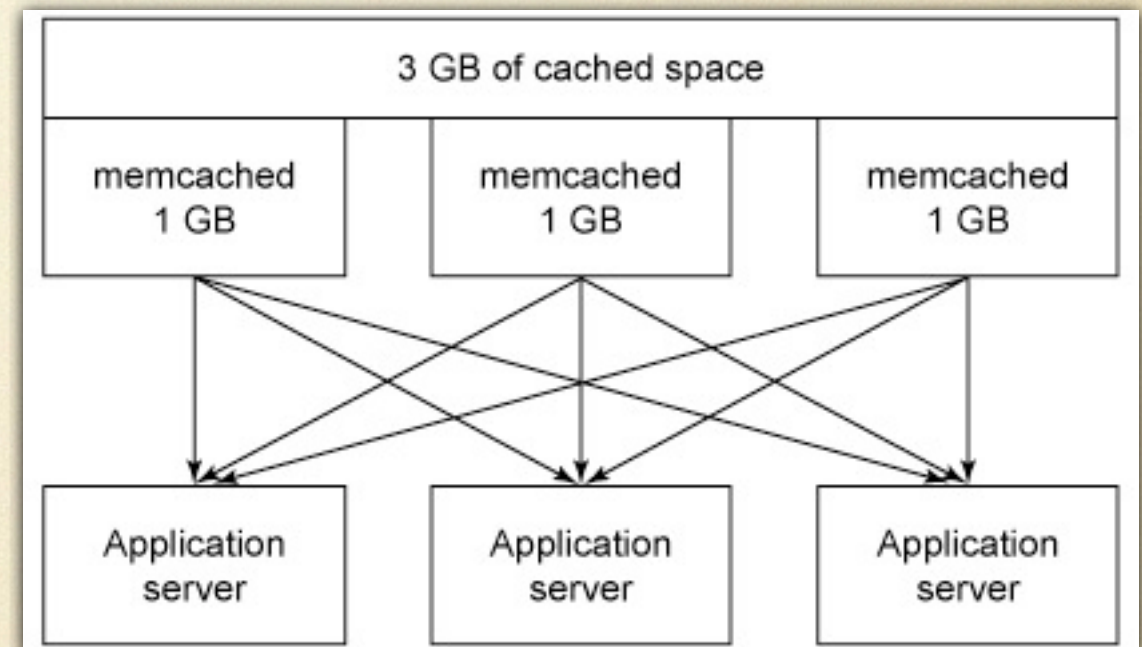
# Abstraction of components



- each service isn't directly offered to users; glueing and wrapping routines will be developed to provide an high level of abstraction
- updates of core services doesn't influence the user applications
- higher flexibility in defining API



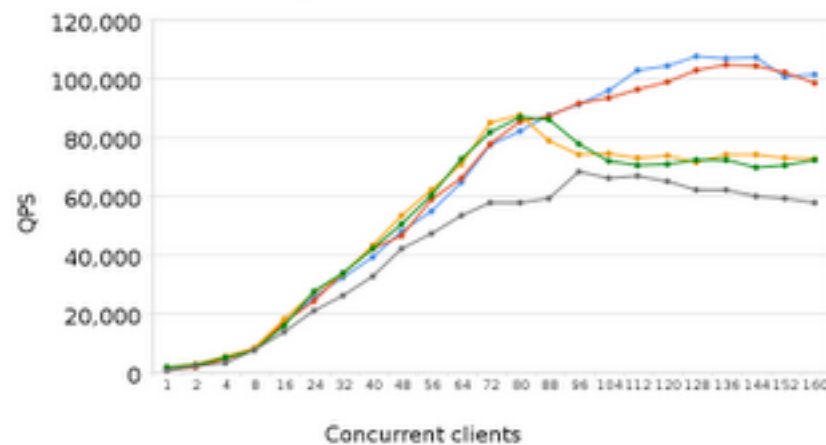
# memcached



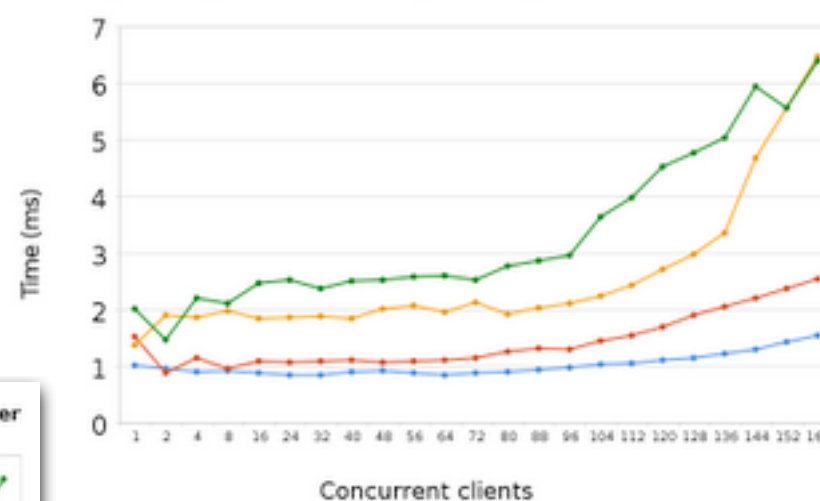


# memcached performance

**Gets by PK - 16 core server**

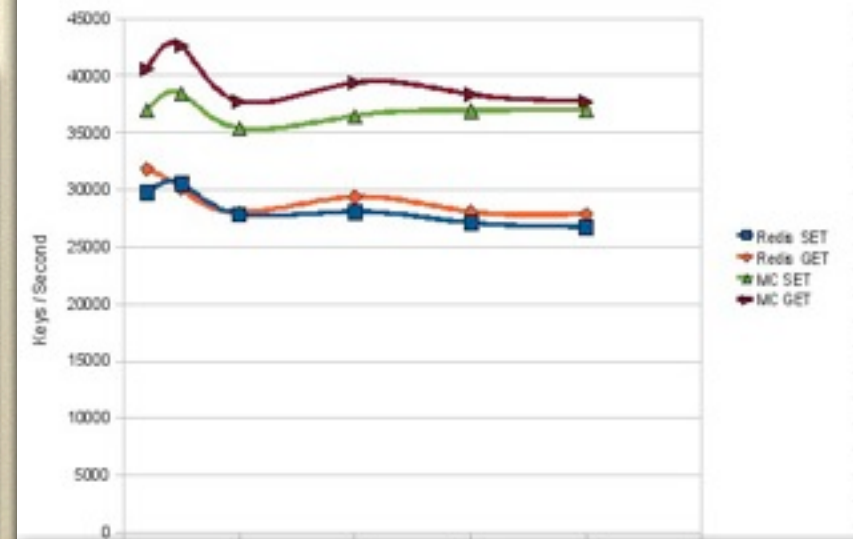


**Gets by PK - response time - average & 98th percentile - 16 core server**

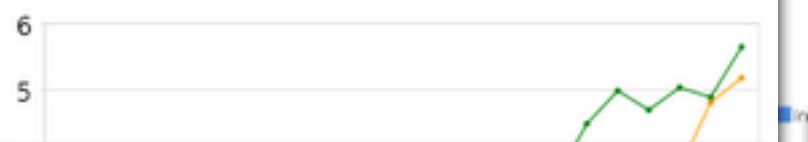


**Redis vs Memcache**

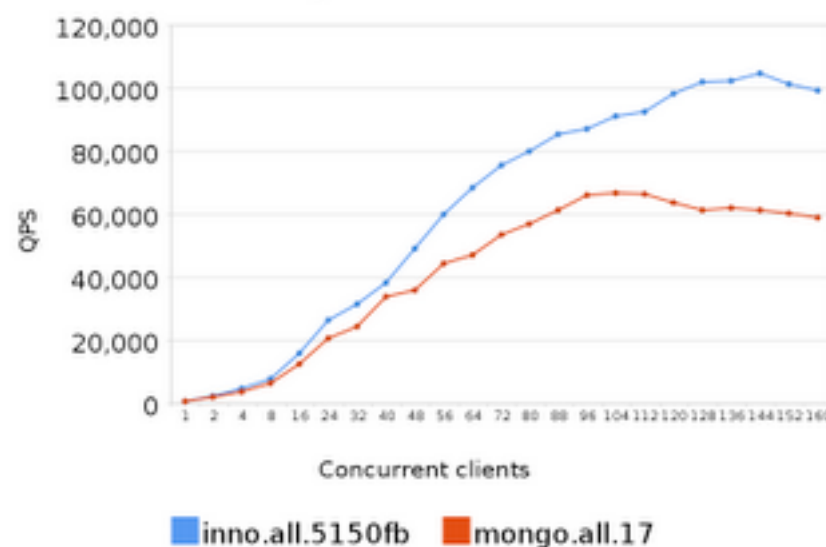
Key Length Bench, 3 byte values



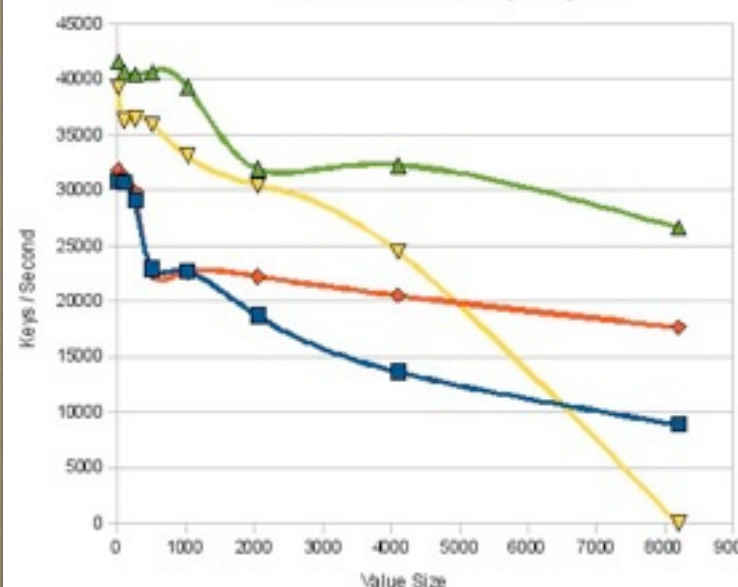
**Gets by PK - response time - average & 98th percentile - 8 core server**



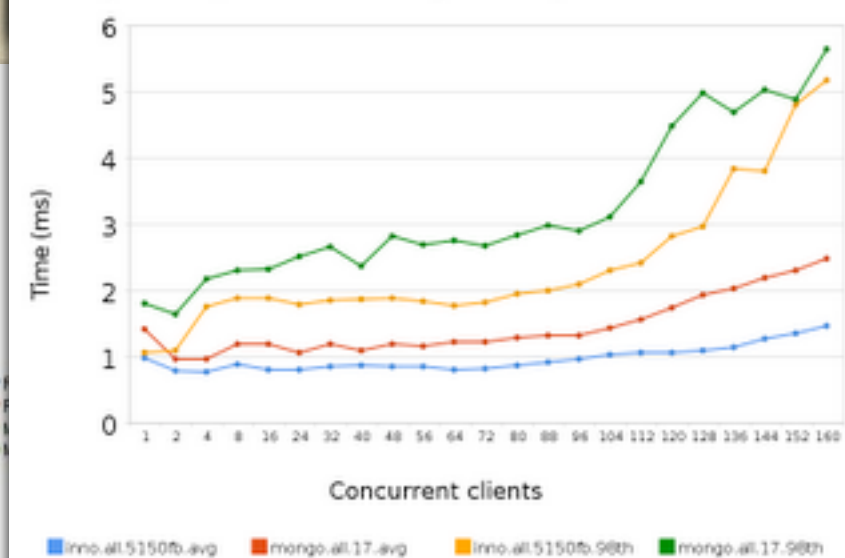
**Gets by PK - 8 core server**



**Redis vs Memcache**  
Value Size Bench, 10 byte key size

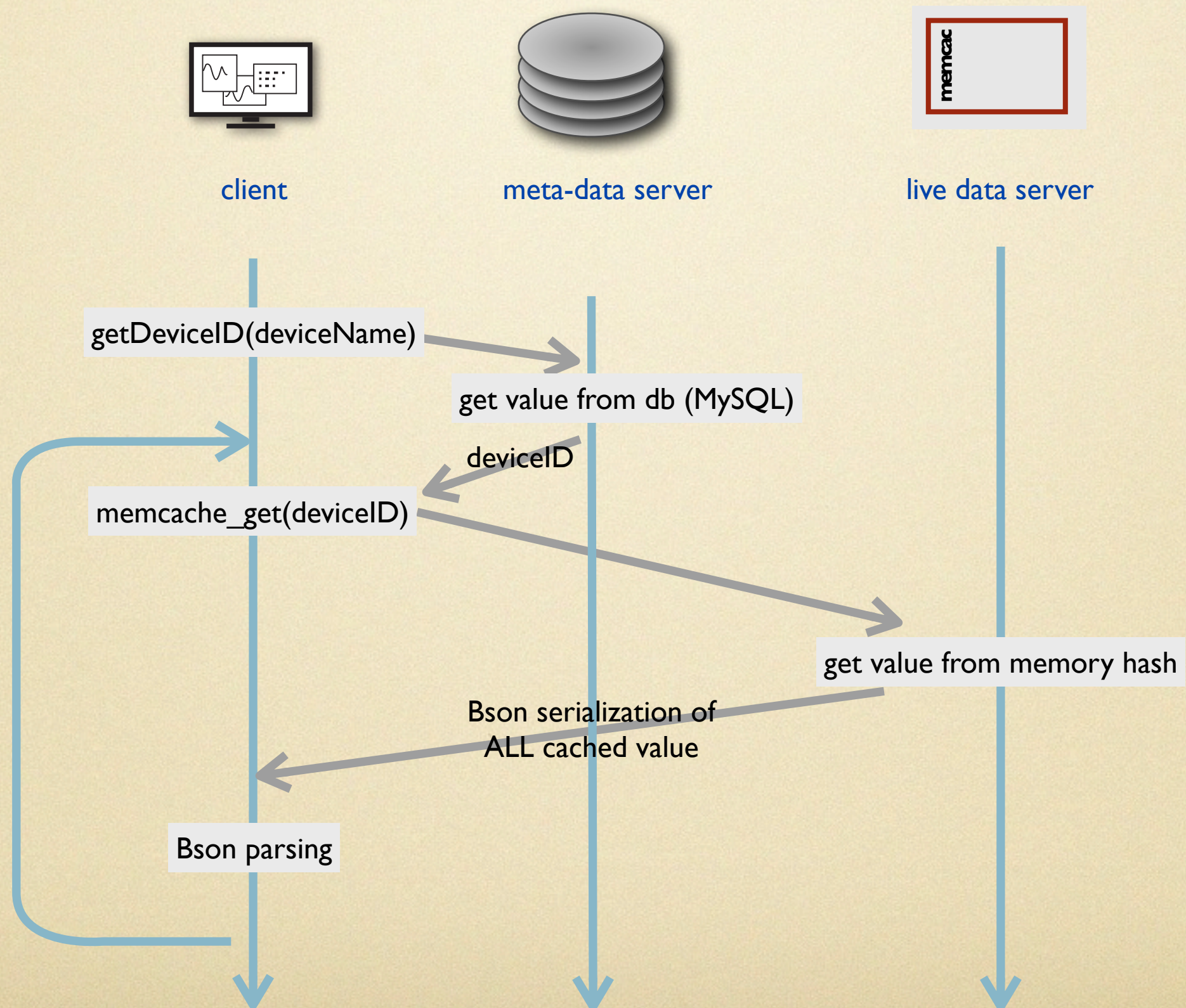


**Gets by PK - response time - average & 98th percentile - 8 core server**



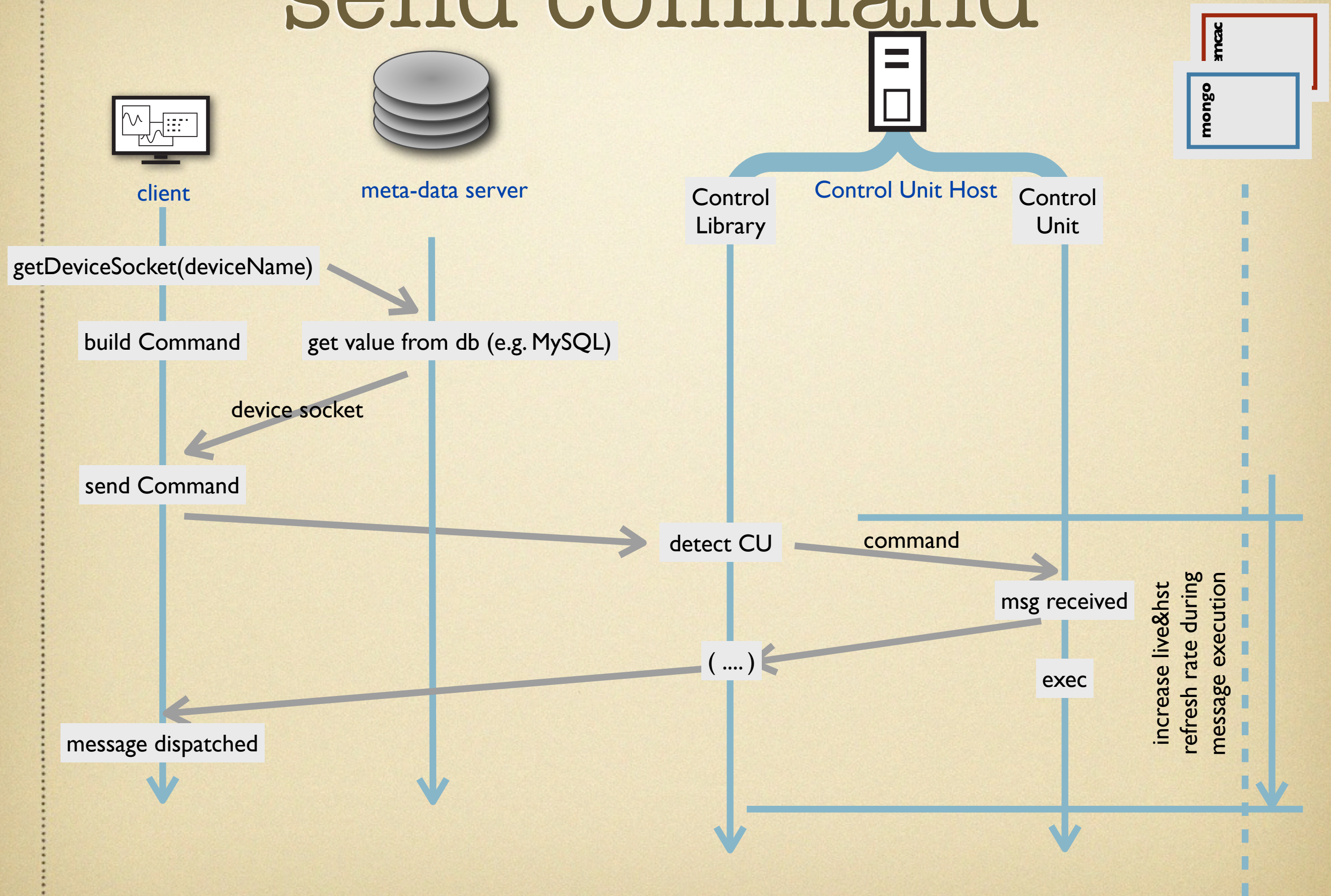


# pull live data



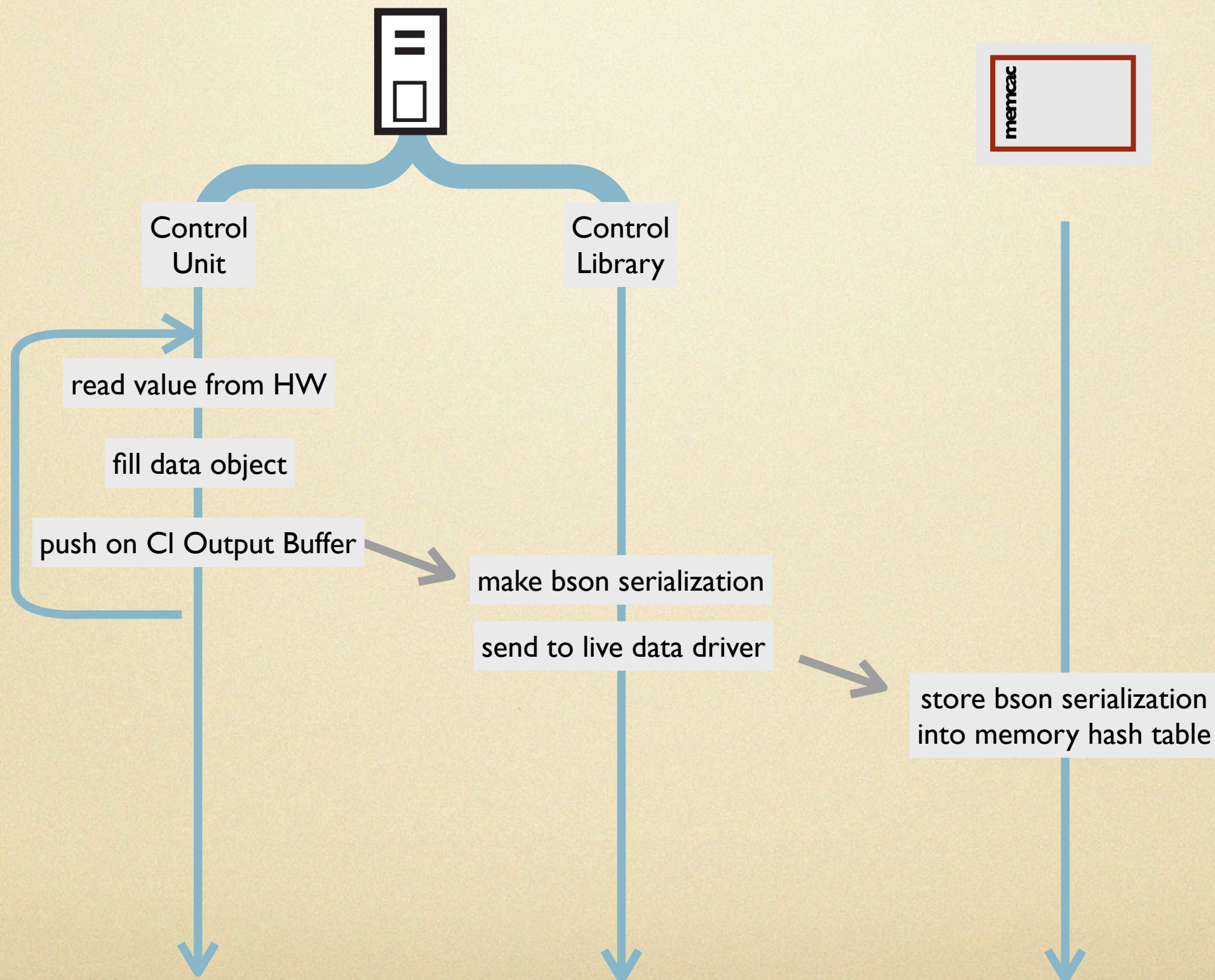


# send command



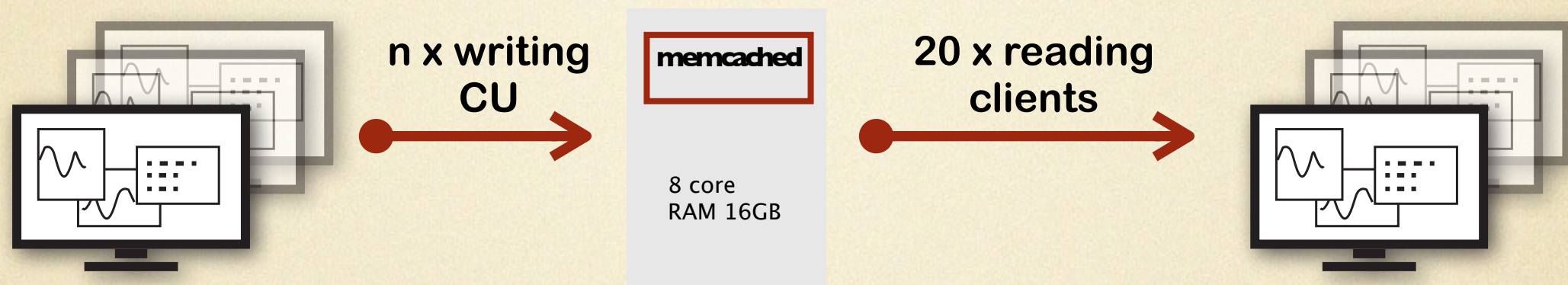


# push live data





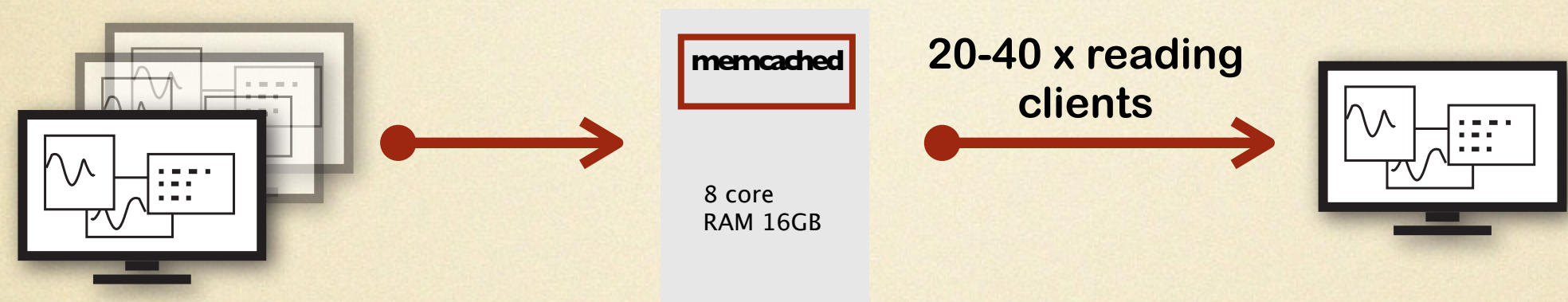
# test#3.1



writing every... (msec)	#CU (Write)	#clients (Read)	#servers	#processes/ server	CPU load (%)
20	60	20	1	1	3-5
20	80	20	1	1	4-6
20	80	20	2	1	2-3
50	60	20	1	1	1-3
50	80	20	2	1	0-2
100	60	20	1	1	?
100	80	20	2	1	?



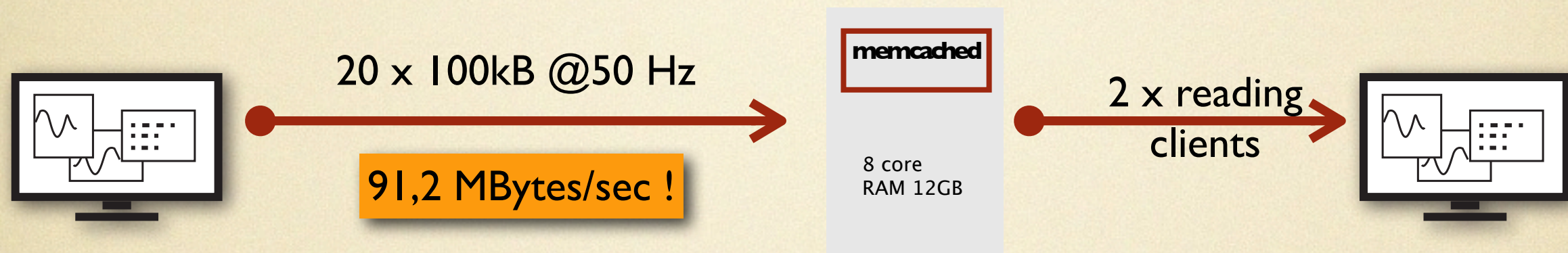
# test#3.2



writing every... (msec)	#CU (Write)	#clients (Read)	#servers	#processes/ server	CPU load (%)
20	80	20	1	4 (1 per core)	2-3
20	80	40	1	4 (1 per core)	2-3
		40	1	4 (1 per core)	0

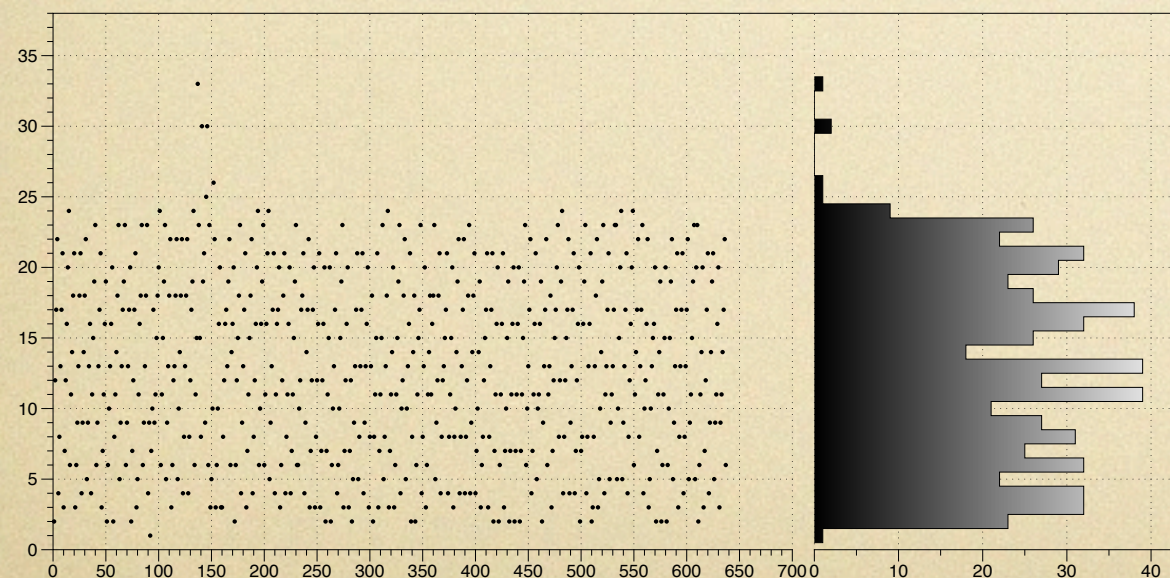


# test#4

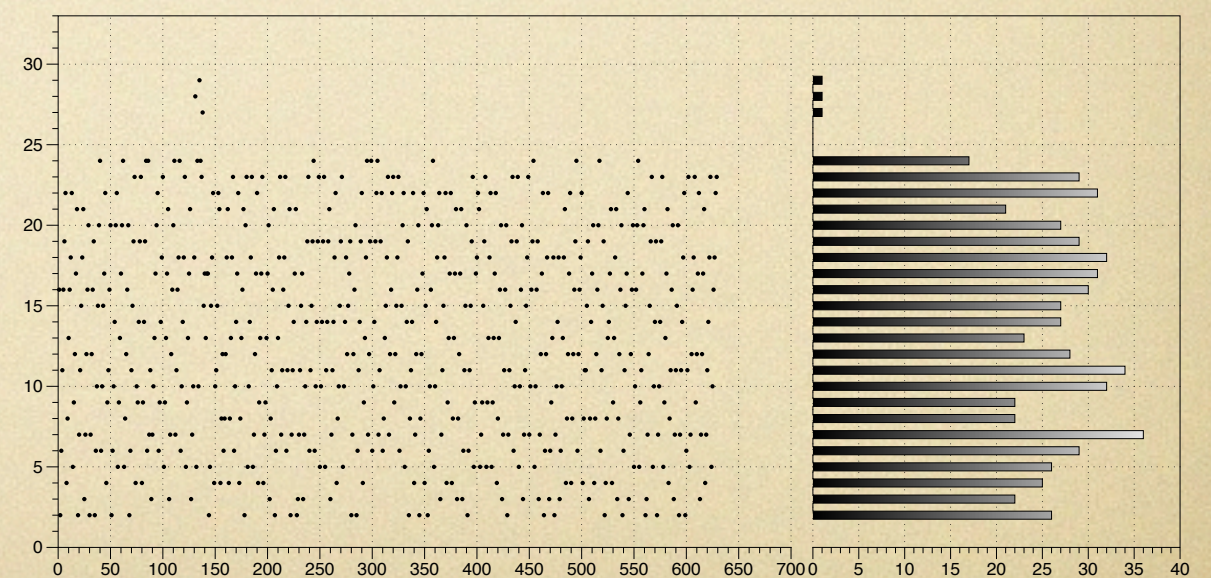


PID	USER	PR	NI	VIRT	RES	SHR	%CPU	MEM	TIME+	COMMAND
28059	dbuser	15	0	72236	10m	616	11.0	0.1	3:50.82	memcached
28066	dbuser	15	0	129m	5688	628	11.0	0.0	3:09.89	memcached
28052	dbuser	15	0	69812	8024	612	7.0	0.0	2:13.86	memcached
28074	dbuser	15	0	67568	5816	616	4.0	0.0	1:29.09	memcached

s4\_hardwareI\_w20\_m20\_buff100000\_rd10.log



s4\_hardwareI\_w20\_m20\_buff100000\_rd12.log





# conclusions and future plans

motivated by the results of preliminary tests and consistency of the overall design:

- continue R&D for completing system design and continue stress tests of components
- prepare a prototype to be tested on the field (test the system during real-life DAFNE & SPARC operations)
- finalize the project as a candidate for the SuperB Control and DAQ System
- evaluate costs, man power and define time schedule