THE O² PROJECT

Giulio Eulisse (CERN)

ALICE IN RUN 2



From O(1) kHz single events...

Pb - Pb



ALICE IN RUN 3

...to 50kHz of continuous readout data.



Overlapping events in TPC with realistic bunch structure @ 50 kHz Pb-Pb Timeframe of 2 ms shown (will be 10-20 ms in production) Tracks of different collisions shown in different colour



CHALLENGES FOR ALICE IN RUN 3

- **Reconstruct 50x more** events online (minbias events only).
- detectors.
- > All of the above in a "flat budget" scenario...

> Store 50x more events (needs TPC compression factor 20x compared to Run 2 raw data size) > Reconstruct TPC data in continuous readout in combination with data coming from triggered

> All of the above while deploying a completely new detector readout and while performing substantial upgrades to the detector itself (new ITS, new GEM for TPC readout, ...).

DESIGNING A NEW COMPUTING ARCHITECTURE FOR ALICE IN RUN 3: ALICE 02

ALICE can cope with the challenges of Run 3 only by a radical redesign of its software and computing architecture.

- > New architecture based on the experience accumulated in the ALICE HLT during Run 1 / Run 2.
- cost for storage.
- > Simplified Data Model in order to improve I/O performance.
- > Appropriately chosen algorithms providing higher throughput for negligible loss of physics
- > Ability to port software components in a gradual manner.
- Close collaboration with the physics community in order to organise analysis efforts.
- > Close collaboration with GSI and FAIR experiments on a common software stack (ALFA).

> Focus on online data compression, only keeping raw data and AOD objects, trading computational

performance. Algorithms tuned for vectorisation / exploitation of hardware accelerators (GPUs).

ALICE IN RUN 3: POINT 2



TIMEFRAME

Data quantum will not be the event, but the "Timeframe".

- ► ~10GB after timeframe building. Vast majority in TPC clusters.
- entropy encoding.
- \succ 50x the number of collisions of Run 2.
- > All MinBias. We need to (lossly) compress information, not filter it.

> ~23ms worth of data taking in continuous readout. Equivalent to 1000 collisions. Atomic unit.

 \blacktriangleright Compressed to $\sim 2GB$ after asynchronous reconstruction, thanks to track-model-compression, storing clusters instead of ADC values, tailored fixed point integer format, logarithmic precision,



SYNCHRONOUS RECONSTRUCTION: GPUS AS FIRST CLASS CITIZENS

Synchronous processing will actually be possible thanks to GPU utilisation for TPC tracking. One modern GPU replaces 40 CPU cores. Changing the algorithm gives an additional 20x - 25x speedup with comparable quality.





David Rohr

Asynchronous Reconstruction

- ► Follows the Data Taking period
- and improved reconstruction software.
- Single persistent analysis object output Analysis Object Data
- \blacktriangleright Processing on O² Facility + T0 (70% of CTF volume), T1 (30% of CTF volume).
- > After 2nd cycle CTF will remain only on tape (removed from the disk buffer) and any subsequent cycle will have to wait until LHC LS.

> 2 processing cycles per data taking year, with increasingly sophisticated calibration



ALICE IN RUN 3: POINT 2



TRANSPORT LEVEL SYSTEM ARCHITECTURE





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ALICE 02: SOFTWARE FRAMEWORK IN ONE SLIDE

Transport Layer: ALFA / FairMQ¹

Standalone processes for deployment flexibility. ► Message passing as a parallelism paradigm. ► Shared memory backend for reduced memory usage and improved performance.

WHY FAIRMQ?

Separation of Concerns: From the arch data transport from the system description.



Performant transport: collaboration with FAIR experiments and GSI allows sharing of highly skilled engineers to work on the performance critical parts related to transport.

Separation of Concerns: From the architectural point of view, it allows ALICE to factor out





ALICE 02: SOFTWARE FRAMEWORK IN ONE SLIDE

Data Layer: 02 Data Model

Transport Layer: ALFA / FairMQ¹

Message passing aware data model. Support for multiple backends: **Simplified, zero-copy** format optimised for performance and direct GPU usage. Useful e.g. for TPC reconstruction on the GPU.

- tools.

ROOT based serialisation. Useful for QA and final results.

> Apache Arrow based. Useful as backend of the analysis ntuples and for integration with other

Standalone processes for deployment flexibility.

► Message passing as a parallelism paradigm.

Shared memory backend for reduced memory usage and improved performance.





O2 DATA MODEL

A timeframe is a collection of (header, payload) pairs. Headers defines the type of data. Different header types can be stacked to store extra metadata (mimicking a Type hierarchy structure). Both header and payloads should be usable in a message passing environment.



Different payloads might have different serialisation strategies. E.g.:

- > QA histograms: serialised ROOT histograms.
- > AOD: some columnar data format. Multiple solutions being investigated.

> TPC clusters / tracks: flat POD data with relative indexes, well suitable for GPU processing.



O2 DATA MODEL



Reconstruction Step 2

Messages being exchanged in O2 have a (header, payload) structure where the header describes the contents of the subsequent payload.

- > Origin represents the Detector or Component that first created the message (e.g. TPC)
- > **Description** is the data type of the payload (e.g. CLUSTERS),
- **Subspecification** can be used to encode extra information (e.g. TPC sectors)
- > Timestamp / Timerange indicates the Timeframe it belongs to.











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DISTRIBUTED SYSTEMS ARE HARD



There are only **two hard problems** in distributed systems:

- 2. Exactly-once delivery
- 1. Guaranteed order of messages
- 2. Exactly-once delivery



DISTRIBUTED SYSTEMS ARE HARD



There are only **two hard problems** in distributed systems:

- 2. Exactly-once delivery
- 1. Guaranteed order of messages
- 2. Exactly-once delivery

Since too many people did not get the joke, we started thinking how to simplify this for the user, as a result we decided to build a **data flow engine (pipelines!)** on top of our distributed system backend.



ALICE 02: SOFTWARE FRAMEWORK IN ONE SLIDE

Data Processing Layer (DPL)

Data Layer: 02 Data Model

Transport Layer: ALFA / FairMQ¹

Abstracts away the hiccups of a distributed system, presenting the user a familiar "Data Flow" system.

Message passing aware data model. Support for multiple backends: **Simplified, zero-copy** format optimised for performance and direct GPU usage. Useful e.g. for TPC reconstruction on the GPU.

- other tools.

Standalone processes for deployment flexibility.

Reactive-like design (push data, don't pull) > Declarative Domain Specific Language for topology configuration (C++17 based).

> Integration with the rest of the production system, e.g. Monitoring, Logging, Control. > *Laptop mode*, including graphical debugging tools.

ROOT based serialisation. Useful for QA and final results.

> Apache Arrow based. Useful as backend of the analysis ntuples and for integration with with

► Message passing as a parallelism paradigm.

Shared memory backend for reduced memory usage and improved performance.









DPL: IMPLICIT WORKFLOW DEFINITION





DPL converts a physics oriented *implicit description of the workflow*





DPL: BUILDING BLOCK

A DataProcessorSpec defines a pipeline stage as a building block.

- Specifies inputs and outputs in terms of the O2 Data Model descriptors.
- Provide an implementation of how to act on the inputs to produce the output.
- Advanced user can express possible data or time parallelism opportunities.







DATA PROCESSING LAYER: IMPLICIT TOPOLOGY





Data Processing Layer

Topology is defined implicitly. Topological sort ensures a viable dataflow is constructed (no cycles!). Laptop users gets immediate feedback through the debug GUI. Service API allows integration with non data flow components (e.g. Control)





inputs/relayed/pending

min timestamp: 0, max timestamp: 1529656515244

- ► A(41498)
- ► B(41499)
- ► C(41500)
- ► D(41501)

	Device Inspector	
	▼ Channels	
	# channels: 2	
	Inputs:	
	Name Port	
	from_B_to_D 22002	
	Name Port	
	▶ Data relayer	
Debug GUI		
🔻 Select metric	▼ Driver information	
	Numer of running devices: 4	
	Plau Pause Step	
	Workflow options:	
in	anInt 1	
in	aFloat 2.000000	
	aDouble 3.000000	
	ascring roo	
	State stack (depth 1)	
	#0: RUNNING	









🔽 Calaat matuia (
Select metric		▼ Drive
	in; in;	Numer o O Play Workflo
		anInt
	in	aFloat
		aDouble
		aString
		aBool
		State s
		40. DU

▼ Driver information				
Numer of running devices: 4				
🔵 Play 🔵 Pause 🔵 Step				
Workflow options:				
anInt	1			
aFloat	2.000000			
aDouble	3.000000			
aString	foo			
aBool	true			
State stack (depth 1)				
#0: RUNNING				



- ► A(64674)
- ► B(64675)
- ► C(64676)
- ► D(64677)

			Device Insp	ector		
	▼ Channels					
			<pre># channels: Inputs: Name from_B_to_D from_C_to_D Outputs: Name ▶ Data rel</pre>	2 ayer	Port 22002 22003 Port	
▼ D(64677)						
				log fil	tan	
				LOG T11	L L I	
				Log sta	and drigge	r
				Log sta	op trigger	
Stop logging	INFO				V	Log
[10:53:30][INFO]	<pre>from_C_to_D[0]:</pre>	in: 0	(0 MB) out: 0	0 (0 MB)		
[10:53:30][INFO]	from_B_to_D[0]:	in: 0.	999001 (0.000	0131868	MB) out: 0	9 (0
[10:53:31][INFO]	from_C_to_D[0]:	in: 0 	(0 MB) out: 0	9 (0 MB) 2 (0 MB)		
[10:53:31][INFO]	from_B_to_D[0]:	in: 0	(0 MB) OUC: ((0 000432 MB)	a (o mb) V outro 0	(0 MB)	
[10:53:32][INFO]	from B to D[0]:	in: I	(0.000102 MD) (0 MB) out: 0	, 00C, 0 A (A MB)	(8 MD)	
[10:53:33][INFO]	from C to D[0]:	in: 0	(0 MB) out: 0) (0 MB)		
[10:53:33][INFO]	from_B_to_D[0]:	in: 1	(0.000132 MB)) out: 0	(0 MB)	
[10:53:34][INFO]	from_C_to_D[0]:	in: 0	(0 MB) out: 0	0 (0 MB)		
[10:53:34][INFO]	<pre>from_B_to_D[0]:</pre>	in: 0	(0 MB) out: 0	0 (0 MB)		
[10:53:35][INFO]	<pre>from_C_to_D[0]:</pre>	in: 0	(0 MB) out: 0	9 (0 MB)		
[10:53:35][INFO]	<pre>from_B_to_D[0]:</pre>	in: 0	(0 MB) out: 0	9 (0 MB)		
[10:53:36][INFO]	from_C_to_D[0]:	in: O	(0 MB) out: 0	9 (0 MB)		
[10:53:36][INFO]	from_B_to_D[0]:	in: 1	(0.000132 MB)) out: 0	(0 MB)	
[10:53:37][INFO]	from_C_to_D[0]:	in: 0.	995025 (0.000	9131343 262697 M	MB) out: (J (0 ∕∩J
		1117-11	99885 (8.888	-02007 W	B) OGC: 0	101

Workflow options:







		Device Inspector	
		▼ Channels	
		# channels: 2 Inputs:	
		Name from_A_to_C Outputs:	Port 22001
		Name from_C_to_D	Port 22003
	▼ Driver inform	nation	
GUI s. ourse nalisa Offline zynek	Numer of runnin Play Pau: Workflow option anInt aFloat aDouble aString aBool State stack (de #0: RUNNING	ng devices: 4 se Step ns: 1 2.000000 3.000000 foo true epth 1)	



```
#include "Framework/runDataProcessing.h"
 3 using namespace o2::framework;
 5 AlgorithmSpec simplePipe(std::string const &what) {
    return AlgorithmSpec{ [what](ProcessingContext& ctx) {
 b
       auto bData = ctx.outputs().make<int>(OutputRef{what}, 1);
    } };
 8
9
10
11 WorkflowSpec defineDataProcessing(ConfigContext const&specs) {
     return WorkflowSpec{
12
    { "A", Inputs{}, {OutputSpec{{"a1"}, "TST", "A1"}, OutputSpec{{"a2"}, "TST", "A2"}},
13
14
       AlgorithmSpec{
15
         [](ProcessingContext &ctx) {
16
          auto aData = ctx.outputs().make<int>(OutputRef{ "a1" }, 1);
          auto bData = ctx.outputs().make<int>(OutputRef{ "a2" }, 1);
17
18
19
20
      "B", {InputSpec{"x", "TST", "A1"}}, {OutputSpec{{"b1"}, "TST", "B1"}}, simplePipe("b1")},
21
     [ "C", {InputSpec{"x", "TST", "A2"}}, {OutputSpec{{"c1"}, "TST", "C1"}}, simplePipe("c1")},
22
      "D", {InputSpec{"b", "TST", "B1"}, InputSpec{"c", "TST", "C1"}}, Outputs{},
23
24
       AlgorithmSpec{[](ProcessingContext &ctx) {}}
25
26
     };
27 }
```

The previous example (GUI included) requires 27 user's SLOC.







exec





<topology id="o2-dataflow">

- <decltask id="A">
- <exe reachable="true">../bin/o2DiamondWorkflow --id A ...</exe>
- <exe reachable="true">../bin/o2DiamondWorkflow --id B ...</exe>
- <decltask id="C">
- <exe reachable="true">../bin/o2DiamondWorkflow --id C ...</exe>
- <decltask id="D">
- <exe reachable="true">../bin/o2DiamondWorkflow --id D ...</exe>





Integration with O2 Control system.





ANALYSIS MODEL: RUN 2

In order to offset the costs of reading data, ALICE has as strong tradition of organised analysis (i.e. trains):

- ► Users provide "wagons", organised in "trains". Trains run on the Grid.
- ► Data is read only once per train, wagons get applied to it.
- > Data is kept in a generic C++ object store, backed by ROOT, as you know.
- Slow sites / site issues is what dominates performance.





Resource Share Projection

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O2 COMPUTING MODEL IN ONE SLIDE



CALIBRATION RECONSTRUCTION STORAGE

ANALYSIS MODEL: RUN 3

Solid foundations: the idea of organised analysis will remain. Improve on the implementation.

- > x100 more collisions compared to present setup
- > Do initial analysis on a fraction of the data on fewer, highly performant, Analysis Facilities.
- ► Full analysis on a reduced set of wagons on the Grid ⇒ Prioritise processing according to physics needs. > Streamline data model, reducing generality and features set to improved speed.
- > Explore different compression strategies (e.g. LZ4, Zstd, custom compression code)
- **Recompute** quantities on the fly rather than storing them. CPU cycles are cheap.
- ► Goal is to have each Analysis Facility go through 5PB of AODs every 12 hours (~100GB/s).





REQUIREMENTS FOR THE AOD FORMAT

AOD's data format will have to play well with AliceO2 message passing, shared memory backed, distributed nature.

- > Zero-{Copy, Serialisation, Adjustments}: we want to be able to reuse data between processes.
- **Growable**: *ability to extend columns on the fly.*
- > **Prunable:** *ability to drop columns on the fly.*
- Skimmable: *ability to select only certain rows.*

Strategy: we are willing to lose some degree of generality for performance.



APACHE ARROW: A POSSIBLE SOLUTION FOR IN-MEMORY COLUMNAR FORMAT

"Cross-language development platform for in-memory columnar data."



Well established. Top-Level Apache project backed by key developers of a number of opensource projects: **Calcite**, Cassandra, Drill, Hadoop, HBase, Ibis, Impala, Kudu, **Pandas**, **Parquet**, Phoenix, **Spark**, and Storm.

Very active. 119 contributors, https://github.com/apache/arrow

O2 design friendly. *message passing / shared memory friendly. Support for zero-copy slicing, filtering.*





APACHE ARROW: A FEW TECHNICAL DETAILS

In memory column oriented storage. Full description https://arrow.apache.org/docs/ memory_layout.html. Data is organized in Tables. Tables are made of Columns. Columns are (<metadata>, Array). An Array is backed by one or multiple Buffers.

Nullable fields. An extra bitmap can optionally be provided to tell if a given slot in a column is occupied.

Nested types. Usual basic types (int, float, ..). It's also possible (via the usual record shredding presented in Google's Dremel paper) to support nested types. E.g. a String is a List < Char > .

No (generic) polymorphism. The type in an array can be nested, but there is no polymorphisms available (can be faked via nullable fields & unions).

Suitable for ALICE analysis needs?



APACHE ARROW: INTEGRATION WITH ROOT

The main concern here is of course "how do I use this from ROOT"?



RDataFrame: new component of ROOT for "declarative analysis". Modular design allows

Initial integration of Arrow with ROOT has already been provided by ALICE and merged by the ROOT team.

Bonus: ROOT gets seamless integration with many OpenSource projects which you can mention to impress your friends and that make your CV look good to head-hunters.



DPL AS AN INTEGRATION PLATFORM FOR 02





O2 Monitoring and InfoLogger integration





WP7 / Quality Control

DataSampling





MID Filtering Chain

TPC Event Displays







TPC reconstruction



BACKUP



```
DataProcessorSpec{
  "A",
  Inputs{
   InputSpec{"a", "TPC", "CLUSTERS"}
  },
  Outputs{
   OutputSpec{{"b"}, "TPC", "TRACKS"}
  <u></u>
  AlgorithmSpec{
    [](ProcessingContext &ctx) {
       auto track = ctx.outputs().make<Track>(OutputRef{ "b" }, 1);
```







```
DataProcessorSpec{
  "A",
  Inputs{
   InputSpec{"a", "TPC", "CLUSTERS"}
  },
  Outputs{
   OutputSpec{{"b"}, "TPC", "TRACKS"}
  ſ,
  AlgorithmSpec{
    [](ProcessingContext &ctx) {
       auto track = ctx.outputs().make<Track>(OutputRef{ "b" }, 1);
```







```
DataProcessorSpec{
  "A",
  Inputs{
   InputSpec{"a", "TPC", "CLUSTERS"}
  },
  Outputs{
   OutputSpec{{"b"}, "TPC", "TRACKS"}
  AlgorithmSpec{
    [](ProcessingContext &ctx) {
       auto track = ctx.outputs().make<Track>(OutputRef{ "b" }, 1);
```







```
DataProcessorSpec{
  "A",
  Inputs{
   InputSpec{"a", "TPC", "CLUSTERS"}
  },
  Outputs{
   OutputSpec{{"b"}, "TPC", "TRACKS"}
  <u>ر</u>
  AlgorithmSpec{
    [](ProcessingContext &ctx) {
       auto track = ctx.outputs().make<Track>(OutputRef{ "b" }, 1);
```





Data is described as pushed through the pipeline.











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