

Motivation
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Tree-level amplitudes
oooooo

Functional reconstruction on \mathbb{Q}
ooo

Finite Fields
oo

BCFW and the maximum-cut
oooooo

The Natural Structure of Scattering Amplitudes

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Università degli studi di Padova

Cortona, 25 May 2018

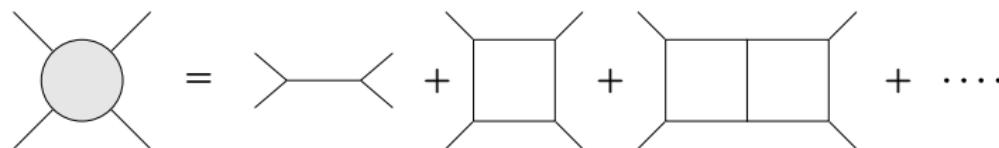
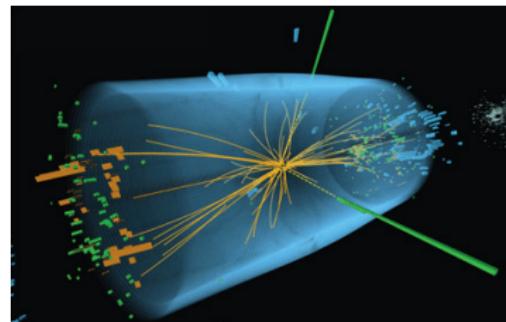
In collaboration with: P. Mastrolia and T. Peraro

Outline

- Motivation
- Numerical tree-level amplitudes from recursive relations
- Functional reconstruction on \mathbb{Q}
- Functional reconstruction on \mathbb{Z}_n
- BCFW tree-level recursion and its relations to the maximum-cut
- Conclusions and outlook

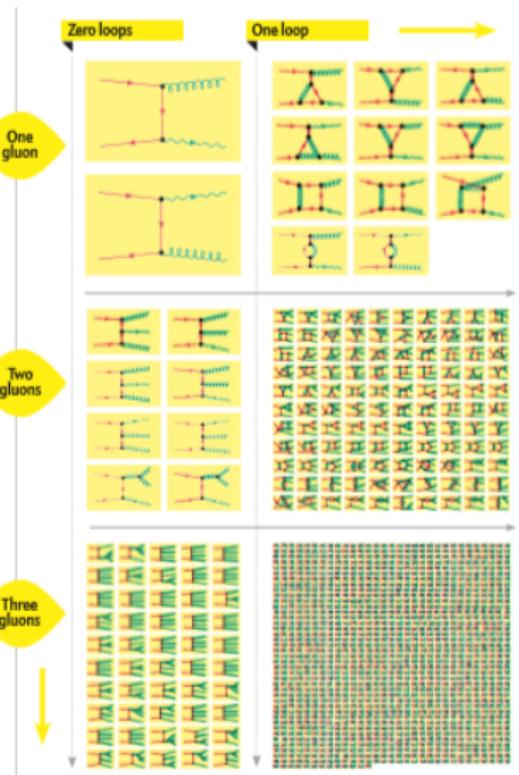
Scattering amplitudes:

- connect theory and experiment
- confirm established models
- open the way for new physics
- infer general structure of a theory



High precision $\xrightarrow{\text{requires}}$ loop corrections

High c.o.m. energy $\xrightarrow{\text{leads to}}$ multi-particle interactions

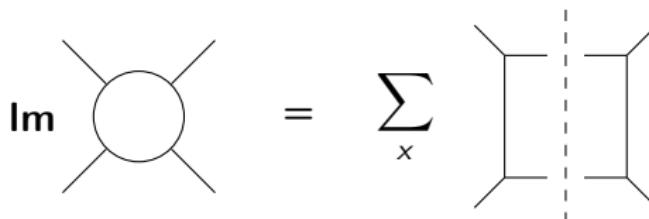


Complexity of the calculations increases quickly with the number of legs and the number of loops.

Efficient techniques needed

- Generalized unitarity $\xrightarrow{\text{extends}}$ optical theorem:

$$\mathbf{M}(i \rightarrow f) - \mathbf{M}^*(i \rightarrow f) = \sum_x \mathbf{M}(i \rightarrow x) \mathbf{M}^*(x \rightarrow f)$$



- unitarity cut:

Tree-level amplitudes $\xrightarrow{\text{building blocks}}$ multi-loop amplitudes

Simple n -guon amplitudes

$$iM(1^\pm, 2^\pm, \dots, n-1^\pm, n^\pm) = 0 \quad [\text{Parke, Taylor '86}]$$

$$iM(1^\mp, 2^\pm, \dots, n-1^\pm, n^\pm) = 0$$

$$iM(1^+, \dots, i^-, \dots, j^-, \dots, n^+) = i \frac{\langle i|j\rangle^4}{\langle 1|2\rangle\langle 2|3\rangle \cdots \langle n-1|n\rangle\langle n|1\rangle}$$

$$iM(1^-, \dots, i^+, \dots, j^+, \dots, n^-) = (-1)^n \frac{[i|j]^4}{[1|2][2|3] \cdots [n-1|n][n|1]}$$



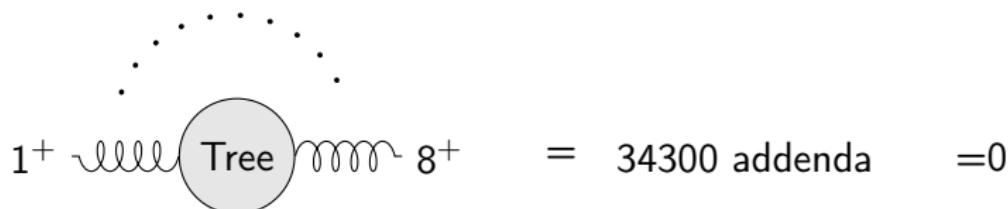
Simple n -guon amplitudes

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Motivation
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○○

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Solution:

Reconstruct the analytic expression of the tree-level amplitudes
from numerical evaluations over finite fields



Large intermediate expressions are replaced by natural numbers.

Important:

We are using functional reconstruction algorithms for rational functions, thus the amplitude must be written as rational function.

N.B. For simplicity, from here on we consider amplitudes involving only gluons.

Spinor-helicity formalism

[Mangano,Parke]

Dirac equation in momentum space for massless particles:

$$i\cancel{p}\psi(x) = 0 \Rightarrow \begin{cases} \cancel{p}U_s(p) = 0 \\ \cancel{p}V_s(p) = 0 \end{cases} \quad s = L, R$$

- Spinors: $U_R(p), U_L(p), \bar{U}_L(p), \bar{U}_R(p)$ \leftrightarrow $|p\rangle, |p], \langle p|, [p|$
- Momenta: $p^\mu = \frac{\langle p|\sigma^\mu|p]}{2}$ $\cancel{p} = |p\rangle[p| + |p]\langle p|$
- Polarizations: $\epsilon_R^\mu(k, r) = \frac{1}{\sqrt{2}} \frac{\langle r|\sigma^\mu|k]}{\langle r|k\rangle}$ $\epsilon_L^\mu(k, r) = \frac{1}{\sqrt{2}} \frac{\langle k|\sigma^\mu|r]}{[k|r]}$
- Invariants: $s_{ij} = (p_i + p_j)^2 = \langle i|j\rangle[j|i]$

Tree-level scattering amplitudes are rational functions of the spinor components.

Amplitude parametrization

$M(p_1, \dots, p_n) \rightarrow M(|p_1\rangle, |p_1], \dots, |p_n\rangle, |p_n])$ rational function

Suitable parametrization of $|p\rangle, |p]$ needed: [Hodges '09; Badger '16; Peraro '15]

momentum-twistor variables

- $3n - 10$ independent variables $\{x_1, \dots, x_{3n-10}\}$

$$3n - 10 = 4n - \underbrace{n}_{\text{onshell}} - \underbrace{4}_{\text{momentum conservation}} - \underbrace{6}_{\text{Lorentz invariance}}$$

- M is rational in the x_i

- x_i allow arbitrary values



momentum conservation and on-shellness always satisfied

■ $n = 5$, a parametrization:

[Badger '16]

$$\begin{pmatrix} |1\rangle & |2\rangle & |3\rangle & |4\rangle & |5\rangle \\ |\mu_1] & |\mu_2] & |\mu_3] & |\mu_4] & |\mu_5] \end{pmatrix} = \begin{pmatrix} 1 & 0 & \frac{1}{x_1} & \frac{1}{x_1} + \frac{1}{x_1 x_2} & \frac{1}{x_1} + \frac{1}{x_1 x_2} + \frac{1}{x_1 x_2 x_3} \\ 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & -1 & -1 & \frac{x_5}{x_4} - 1 \\ 0 & 0 & 0 & \frac{x_4}{x_2} & 1 \end{pmatrix}$$

$s_{12} = x_1$

$s_{23} = x_1 x_4$

$s_{34} = \frac{x_1((x_5 - 1)x_2 x_3 + (x_3 + 1)x_4)}{x_2}$

$s_{45} = x_1 x_5$

$s_{51} = x_1 x_3 (x_2 - x_4 + x_5)$

■ $n = 4$, an amplitude:

$$M_4(1^+, 2^-, 3^+, 4^-) = -\frac{8(1+x_2)^4}{x_1^2 x_2^3}$$

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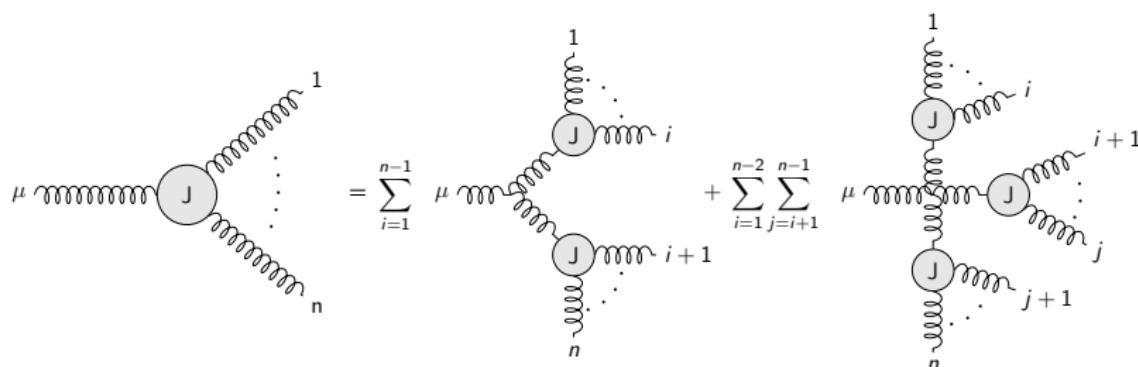
- $n = 4$, an amplitude:

$$M_4(1^+, 2^-, 3^+, 4^-) = -\frac{8(1+x_2)^4}{x_1^2 x_2^3}$$

Off-shell recurrence relation

[Berends, Giele; Mangano,Parke]

$$\mathbf{M}(1, \dots, n+1) \mapsto J^\mu(1, \dots, n)$$



$$\mathbf{M}_{n+1}(1, \dots, n+1) = \lim_{P_{1,n} \rightarrow 0} \epsilon_\mu(n+1) J^\mu(1, \dots, n) P_{1,n}^2$$

On-shell recurrence relation

Complex momenta + residue theorem = Amplitude factorization

$$p_i, p_j \rightarrow p_i(z), p_j(z) \quad z \in \mathbb{C} \quad \Rightarrow \quad M \rightarrow M(z)$$

$$\frac{1}{2\pi i} \oint_{C_\infty} \frac{iM(z)dz}{z} = \sum_{\text{all poles } \alpha} \text{Res}_{z=z_\alpha} \frac{iM(z)}{z} = 0$$



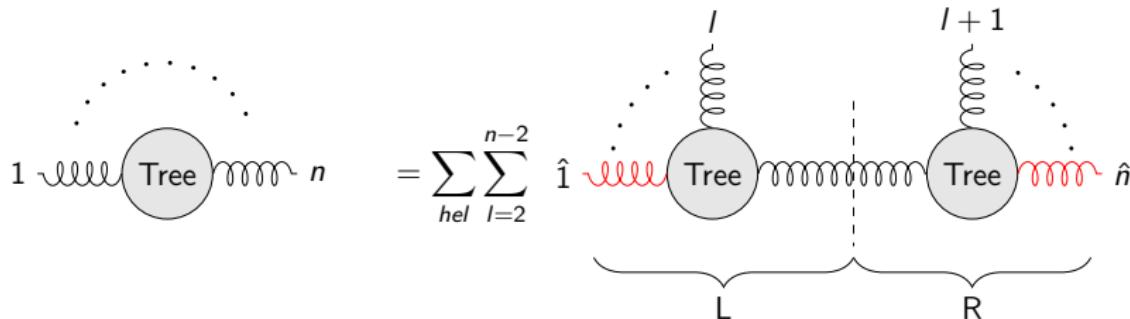
[Britto, Cachazo, Feng, Witten '05]

On-shell recurrence relation

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On-shell recurrence relation

$$\begin{aligned}
 &= \sum_{\text{hel}} \sum_{l=2}^{n-2} \sum_{k=2}^{l-1} \sum_{j=l+1}^{n-2} \text{Tree}_k \text{Tree}_{k+1} \text{Tree}_j \text{Tree}_{j+1} \\
 &\quad \underbrace{\qquad\qquad\qquad}_{L} \qquad \underbrace{\qquad\qquad\qquad}_{R} \qquad \hat{i} \qquad \qquad \hat{n} \\
 &= \sum_{\text{hel}} \sum_{\sigma} S(\sigma) \text{Tree}_2 \text{Tree}_{\sigma(1)} \text{Tree}_{\sigma(2)} \dots \text{Tree}_{\sigma(n-4)} \text{Tree}_{\sigma(n-3)} \text{Tree}_{\sigma(n-2)} \\
 &\quad \hat{i} \qquad \qquad \qquad \qquad \qquad \qquad \qquad \qquad \hat{n} \\
 \text{Building blocks: } & \quad \begin{array}{c} 2^- \\ \text{Tree}_1 \end{array} \quad = \frac{\langle 1|2\rangle^4}{\langle 1|2\rangle\langle 2|3\rangle\langle 3|1\rangle} \quad \begin{array}{c} 2^+ \\ \text{Tree}_1 \end{array} \quad = -\frac{[1|2]^4}{[1|2][2|3][3|1]} \\
 & \quad 1^- \text{Tree}_2 \text{Tree}_3 \text{Tree}_3^+ \quad 1+ \text{Tree}_2 \text{Tree}_3 \text{Tree}_3^- \quad 1-
 \end{array}$$

On-shell recurrence relation

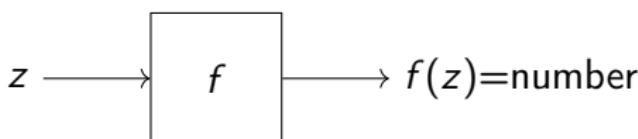
$$\begin{aligned}
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 &\quad \underbrace{\qquad\qquad\qquad}_{1^-} \qquad\qquad\qquad \underbrace{\qquad\qquad\qquad}_{2^+} \qquad\qquad\qquad \underbrace{\qquad\qquad\qquad}_{3^+} \qquad\qquad\qquad \underbrace{\qquad\qquad\qquad}_{1^+} \qquad\qquad\qquad \underbrace{\qquad\qquad\qquad}_{2^+} \qquad\qquad\qquad \underbrace{\qquad\qquad\qquad}_{3^-} \\
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 \end{aligned}$$

On-shell recurrence relation

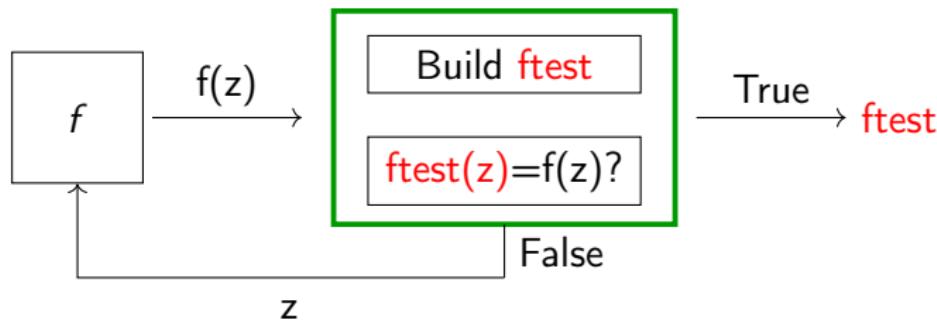
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 \text{Building blocks: } & \quad 1^- \text{Tree}_{2-} \text{Tree}_{3+} = \frac{\langle 1|2\rangle^4}{\langle 1|2\rangle\langle 2|3\rangle\langle 3|1\rangle} \quad 1^+ \text{Tree}_{2+} \text{Tree}_{3-} = -\frac{[1|2]^4}{[1|2][2|3][3|1]}
 \end{aligned}$$

Black-box interpolation

Black-box algorithm:



Reconstruction of the **analytic expression** of f :



Univariate polynomials and rational functions

- Newton's polynomial form:

$$f(z) = a_0 + (z - y_0) \left(a_1 + (z - y_1) \left(\cdots + (z - y_{r-1}) a_r \right) \right)$$

- Thiele's interpolation formula for rational functions:

$$f(z) = a_0 + \frac{z - y_0}{a_1 + \frac{z - y_1}{a_2 + \frac{z - y_2}{\vdots}}}$$
$$\frac{\vdots}{a_{N-1} + \frac{z - y_{N-1}}{a_N}}$$

Example, a polynomial:

- 1 $f(y_0) \xrightarrow{\text{Yields}} a_0 \Rightarrow \text{ftest}(x) = a_0$
- 2 $f(y_1) \xrightarrow{\text{Yields}} a_1 \Rightarrow \text{ftest}(x) = a_0 + (x - y_0)a_1$
- 3 ...

In[67]:= rationalrec[F, z, 4]

$$\begin{aligned}
 \text{Out}[67]= & 3 + (-1 + z) / \\
 & \left(-\frac{741}{580} + (-3 + z) / \left(-\frac{17400}{1637} + (-4 + z) / \left(\frac{140600293}{657044880} + (-5 + z) / \left(\frac{2467158210960}{327086106887} + \right. \right. \right. \right. \right. \\
 & (-6 + z) / \left(\frac{615564163571358245}{5995432842109296} + (-7 + z) / \left(\frac{128795531544061251840}{1474536242413473730979} + \right. \right. \right. \right. \right. \\
 & (-8 + z) / \left(-\frac{30717868632058486438134999}{3713653815285726138460} + \right. \right. \right. \right. \right. \\
 & (-9 + z) / \left(-\frac{119452926980939870886866}{120679006462873659811540965} + \right. \right. \right. \right. \right. \\
 & (-10 + z) / \left(\frac{638117589910825534294864815}{217541179828791209956} + \right. \right. \right. \right. \right. \\
 & (-11 + z) / \left(\frac{17994283078074524}{5602084863015834085935} + \right. \right. \right. \right. \right. \\
 & \left. \left. \left. \left. \left. \left. \frac{-12 + z}{-\frac{124839340988342589}{1550113748} - \frac{342578411}{71} (-13 + z)} \right) \right) \right) \right) \right) \right) \right)
 \end{aligned}$$

In[68]:= % // Simplify

$$\text{Out}[68]= \frac{1 + z^2 + 13 z^4}{3 + z^2 + z^6}$$

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Many rationals are involved in the computation



they often present numerators and denominators described by a huge number of digits



extensive use of arbitrary-precision arithmetic is needed



the computation is strongly slowed down

From rational to natural numbers and back

Finite Field:

$$\mathbb{Z}_n = \{0, \dots, n-1\} \quad n = \text{prime}$$

[Wang '82, Peraro '16]

- From \mathbb{Z} to \mathbb{Z}_n : $\mathbb{Z} \ni a = b \pmod{n} \Leftrightarrow a = b + mn$

Example: $15 = 1 \pmod{7}$ since $15 = 1 + 2 \times 7$

- From \mathbb{Q} to \mathbb{Z}_n : $q = \frac{a}{b} \mapsto (ax(b^{-1} \pmod{n})) \pmod{n}$
- From \mathbb{Z}_n to \mathbb{Q} :

$$\frac{a}{b} = c \pmod{n}$$

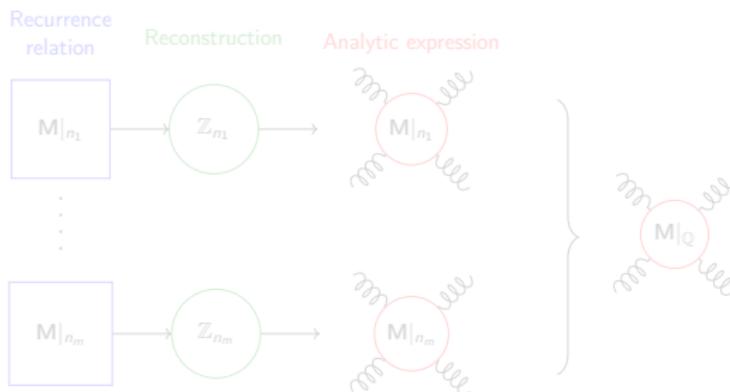
Conditions	Possible pairs (a, b)
none	Infinitely many
$a, b < n$	Finitely many
$a^2, b^2 < \frac{n}{2}$	Only one

$\longrightarrow n$ big enough \Rightarrow unique inverse
 \downarrow
Machine-size limit!

Scattering amplitudes on finite fields

Chinese remainder theorem:

$$\begin{cases} X \equiv x_1 \pmod{n_1} \\ \vdots \\ X \equiv x_m \pmod{n_m} \end{cases} \xrightarrow{\text{Combine into}} X = \tilde{X} \pmod{n_1 \cdots n_m}$$

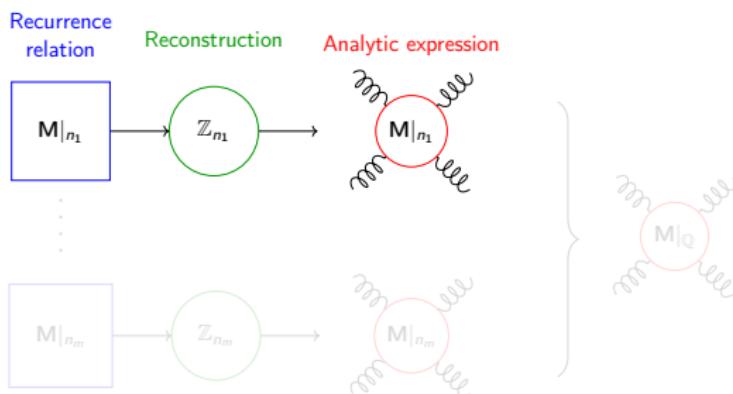


- Off-shell Berends-Giele recurrence [Peraro '16]
- On-shell BCFW recurrence [A.M.]

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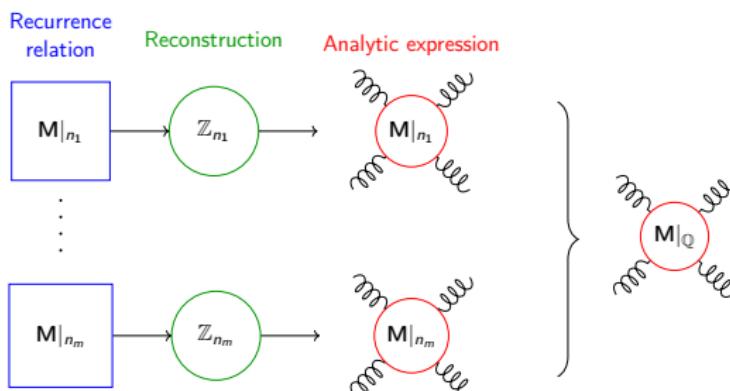


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Scattering amplitudes on finite fields

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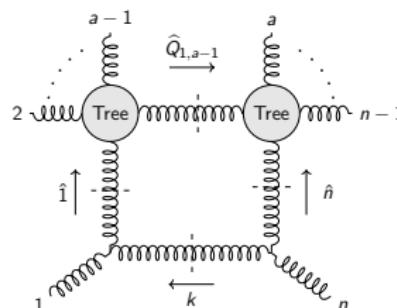
Maximum-cut

- 1 unitarity-cut = one constraint [Mastrolia, Mirabella, Ossola, Peraro '12]
- ℓ loops $\Rightarrow 4\ell$ parameters \Rightarrow at most 4ℓ constraints

Maximum-cut: $D_1 = D_2 = \dots = D_{4\ell} = 0$

Ex: 1-loop \rightarrow quadruple cut:

$$k = x_1 p_1 + x_2 p_4 + x_3 \epsilon_{14} + x_4 \epsilon_{41} \quad \text{span}(\{p_1, p_4\}) \perp \text{span}(\{\epsilon_{14}, \epsilon_{41}\})$$

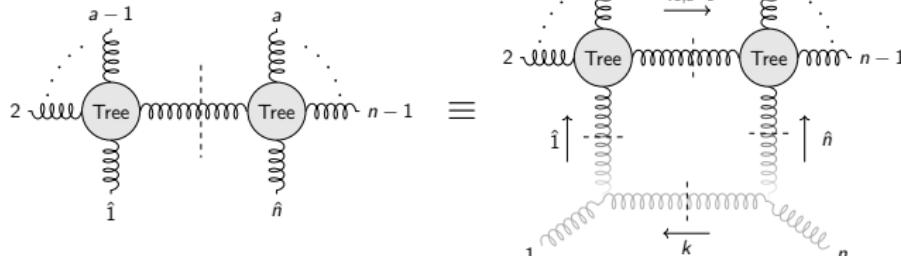


$$\begin{aligned} x_1 &= 0 \\ x_2 &= 0 \\ x_3 &= -\frac{Q_{1,a-1}^2}{Q_{1,a-1} \cdot \epsilon_{1n}} \wedge x_4 = 0 \\ x_3 &= 0 \wedge x_4 = -\frac{Q_{1,a-1}^2}{Q_{1,a-1} \cdot \epsilon_{n1}} \end{aligned}$$

[Britto, Cachazo, Feng '04]

BCFW vs maximum-cut

[Britto, Cachazo, Feng '04; Mastrolia]

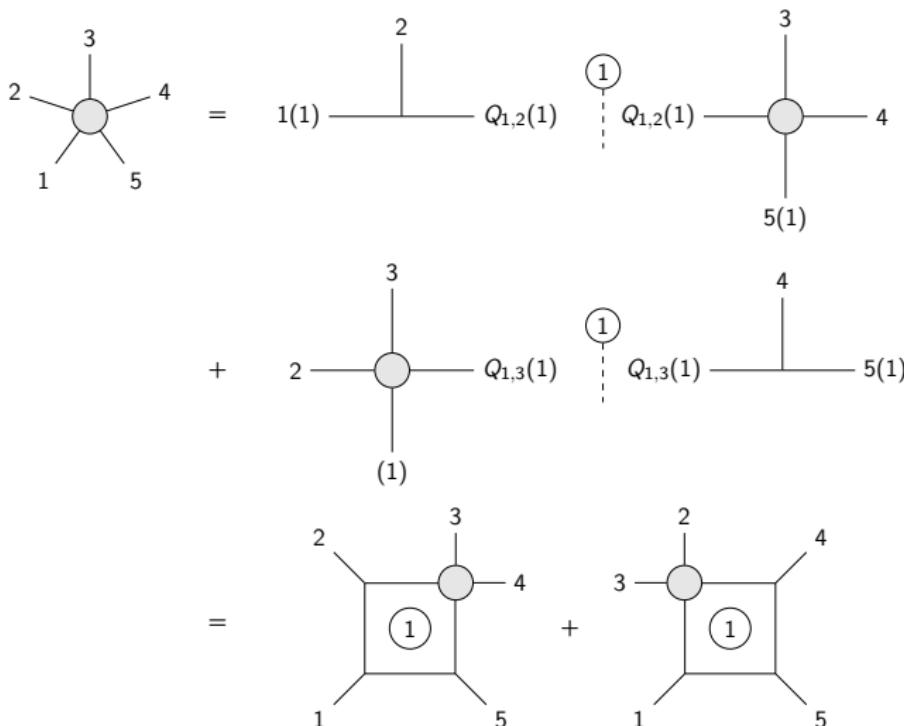


Conditions defining BCFW pole	=	Conditions defining the 4-ple cut
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Systematic extension to higher-loop [A.M.]:

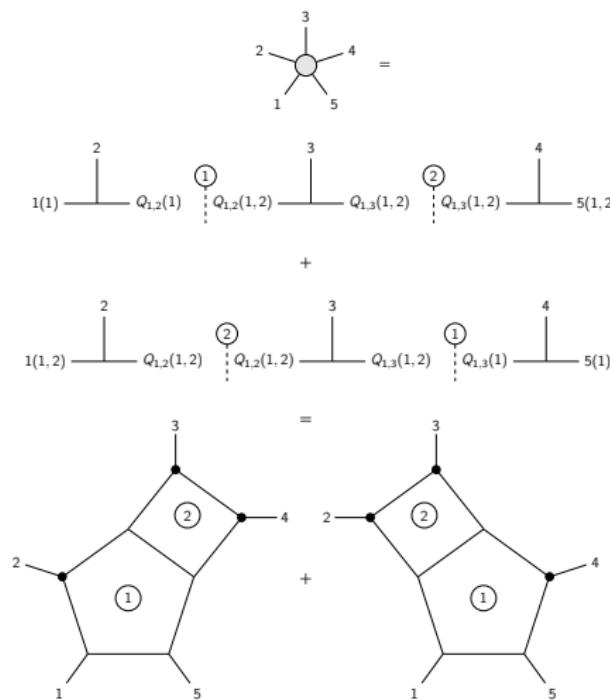
n -gluon tree-level amplitude	$\xrightarrow{\text{BCFW recursion}}$	$(n-3)$ -loop maximum-cut
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5-gluons vs 2-loop



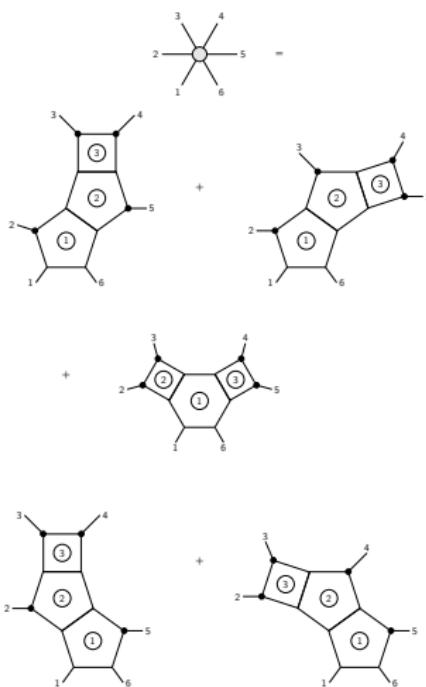
Straight lines still represent gluons and all propagators are considered to be cut.

5-gluons vs 2-loop



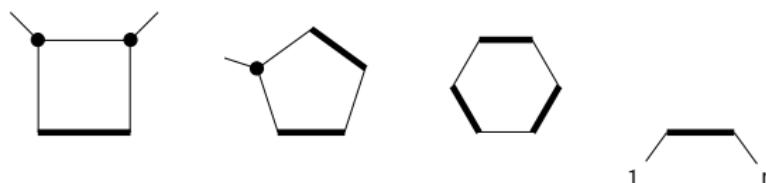
Straight lines still represent gluons and all propagators are considered to be cut.

6-gluons vs 3-loop



Multi-loop building blocks

Maximum-cut graphs from polygons:



Example, $n = 4$:

$$\text{square} \otimes \text{square} = \text{cube}$$

The diagram illustrates the construction of a 3D cube from two 2D squares. On the left, a square and a pentagon are shown with a tensor product symbol (\otimes). To the right, an equals sign ($=$) is followed by a cube where the vertices are labeled 1 through 4. Vertex 1 is at the bottom-front, vertex 2 is at the top-front, vertex 3 is at the top-back, and vertex 4 is at the bottom-back.

Multi-loop building blocks

The final-factorization diagrams can be obtained directly according to the following rules:

- 1 Compute the number of iterations needed for complete factorization: $T = n - 3$
- 2 Build sets of T polygons among those represented in the previous slide that satisfy $N = 5T - 1$, with N sum of all the edges of the figures in the set. Each set must contain at least one square.
- 3 For each set combine the polygons among them and with the base line in all possible ways, considering that:
 - The polygons are to be connected to one another only on the thick lines
 - Polygons with the same number of edges are to be considered indistinguishable
 - Only connected diagrams are admitted

Conclusions

Summary

Any function which can be implemented as a sequence of rational operations is suited for rational reconstruction over finite fields.
Example: tree-level scattering amplitudes.

Outlook

- Implementation of multivariate polynomial/rational reconstruction algorithm.
- Extension to theories with fermions and scalars.
- Extension to massive particles.
- Application of the rational reconstruction to other techniques
- Thorough study of the relation between tree-level BCFW recurrence and multi-loop diagrams.