



GPGPU Evaluation – First experiences in Napoli

Silvio Pardi







Goal of our preliminary tests

•Achieve know-how on GPGPU architectures in order to test the versatility and investigate the possible adoption for some specific tasks interesting for SuperB.



• General Purpose



many-core technology

- Hundred of simple Processing Units
- Designed to match the SIMD paradigm (Single Instruction Multiple Data)







The Hardware Available in Napoli





1U rack NVIDIA Tesla S2050

- ➤ 4 GPU Fermi
- Memory for GPU: 3.0 GB
- Core for GPU: 448
- Processor core clock: 1.15 GHz

2U rack Dell PowerEdge R510

- ➢Intel Xeon E5506 eightcore @ 2.13 GHz
- ▶32.0 GB DDR3
- ➢ 8 hard disk SATA (7200 rpm), 500 GB









- •Cuda compilation tools, release 4.0, V0.2.1221
- •NVIDIA-Linux-x86_64-270.41.34.run
- •cudatoolkit_4.0.17_linux_64_rhel5.5.run
- •gpucomputingsdk_4.0.17_linux.run







Tuning for lazily modality removal

The firsts test on the Tesla S2050 GPU with CUDA C show a starting overhead <u>~2 second</u> due the "context" creation needed by the CUDA toolkit. The context is create **on demand** (**lazily**) and de-allocated when is not used.

SOLUTIONS (Suggested by Davide Rossetti):

- 1. Before the CUDA 4, Create a *dummy process always active that keeps alive the CUDA context* "**nvidia-smi -l 60 -f /tmp/bu.log**"
- 2. Since CUDA ver. 4 use the *-pm* (*persistence-mode*) option from the nvidia-smi command to activate the GPU context.







CUDA C



IBRID code: combination of standard C sequential code with parallel kernel in **CUDA C.**

Each code needs the following main steps:

- ➤ GPU Memory Allocation (CudaMalloc)
- > Data transfer between CPU and GPU (CudaMemCopy H2D)
- > CUDA Kernel execution
- Data transfer between GPU and CPU (CudaMemCopy D2H)







CUDAMALLOC Benchmark









CUDAMEMCPY Benchmark





SuperB Exercise: B-meson reconstruction like algorithm



stage	particelle
1	tracks, K_S , γ , π^0
2	$D^{\pm}_{(s)}, D^0,$ e J/ψ
3	$D_{(s)}^{*\pm}$ e D^{*0}
4	B^{\pm} e B^{0}

SuperB SuperB

Combinatorial problem

Problem Modellization: given N

quadrivectors (spatial components and energy), combine all the couple without Repetition. Then calculate the mass of the new particle and check if the mass is in a range given by input.

GOAL: Understand the impact, benefits and limits of using the GPGPU architecture for this use case, through the help of a toy-model, in order to isolate part of the computation.

The algorithm has been implemented in <u>C CUDA</u>





SuperB SuperB



Cuda implementation (1/2)







Cuda Kernel (2/2)





Super B Super B







Some Ideas

The first experience suggest to continue the investigation expecially in the following ways:

Tuning the memory management: Investigate the possibility to Overlapping Data Transfers and Computation through Async and Stream APIs.

Tuning the Grid/Block/Threads topology

Consider to rearrange the algorithms in term of operations per threads.







Conclusion

In Napoli we started to test the NVIDIA GPGPU architecture and we are investigating how to port HEP code on these architectures.

A first experience using a toy algorithm has show several aspects to take in account in GPU programming:

- •Overhead management
- Memory management
- •Algorithm re-engineerization
- •Work distribution for each Thread
- •Block and Thread topology definition

A lot of work is still due in order to achieve a full understanding of the architecture and the real benefits achievable. There are several work in progress.







END







Comparison malloc e cudaMalloc







Algorithm for delegate a couple of vector to each Thread



```
. . .
//Calcolo del thread ID
sh=threadIdx.x+blockDim.x*threadIdx.y;
tid = blockDim.x*blockDim.y*(blockIdx.x + gridDim.x* blockIdx.y) + sh;
//Algoritmo per il calcolo degli indici giusti in base al thread
if (tid<N2) {
        linIdx=N2-tid:
        i=int(N - 0.5 - sqrt(0.25 - 2 * (1 - linIdx)));
        z=(N+N-1-i)*i;
        j=tid - z/2 + 1 + i;
        if (i==j) {
                i=i-1;
                j=N-1;
        ł
        //Somma di due quadrivettori
        Candidato.x=Gamma[i].x+Gamma[j].x;
        Candidato.v=Gamma[i].v+Gamma[j].v;
        Candidato.z=Gamma[i].z+Gamma[j].z;
        Candidato.Ene=Gamma[i].Ene+Gamma[j].Ene;
        . . .
```



Page-Locked Data Transfers



A inter

cudaMallocHost() allows allocation of pagelocked ("pinned") host memory

- Enables highest cudaMemcpy performance
 - 3.2 GB/s on PCI-e x16 Gen1
 - 5.2 GB/s on PCI-e x16 Gen2

See the "bandwidthTest" CUDA SDK sample

Use with caution!!

Allocating too much page-locked memory can reduce overall system performance



Test your systems and apps to learn their limits

INVIDIA Corporation 2009

-23



